

A Lock-Free concurrent algorithm for Linked lists

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Paper: Mikhail Fomitchev and Eric Ruppert, Lock-Free Linked Lists and Skip Lists. *PODC'04*, 2004.

Plan

1. Introduction
2. The problem
3. The algorithm

1. Introduction

The goals

Simply allow to make concurrent operations on a same shared ordered linked list...but without using OS synchronisation primitives (lock-free).

Operations: **Search** (for a key), **Insert** and **Delete** an element

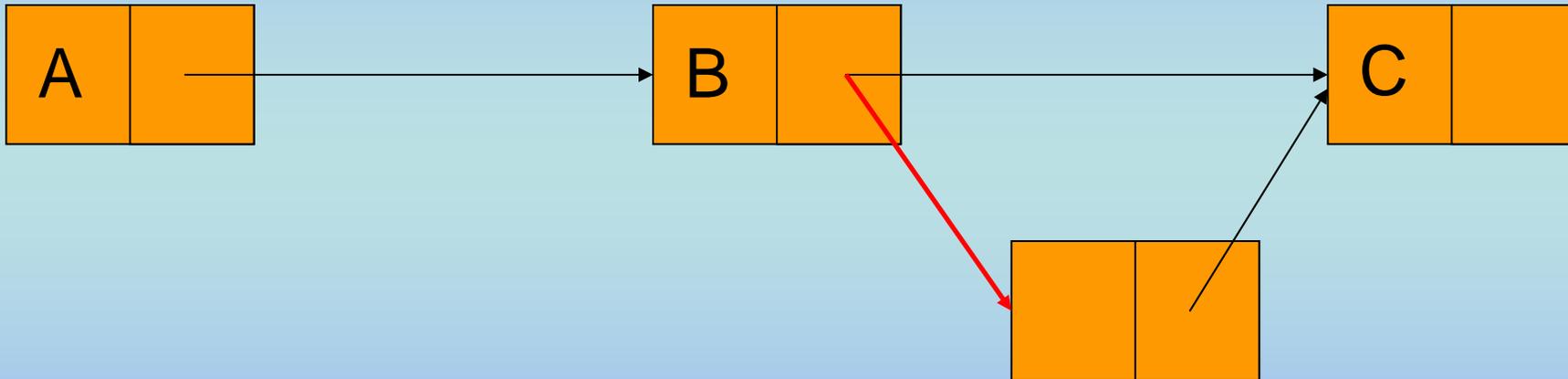
The motivation

- Interested in the CAS instruction
- Already Worked on a Spin-locks algorithm, wanted to see further : **free-locks algorithm**
- Linked lists are everywhere and are the base of all the data structures

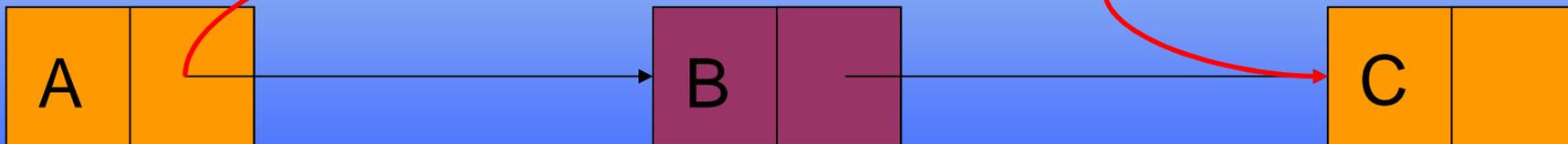
2. The Problem

Reminder

Insertion



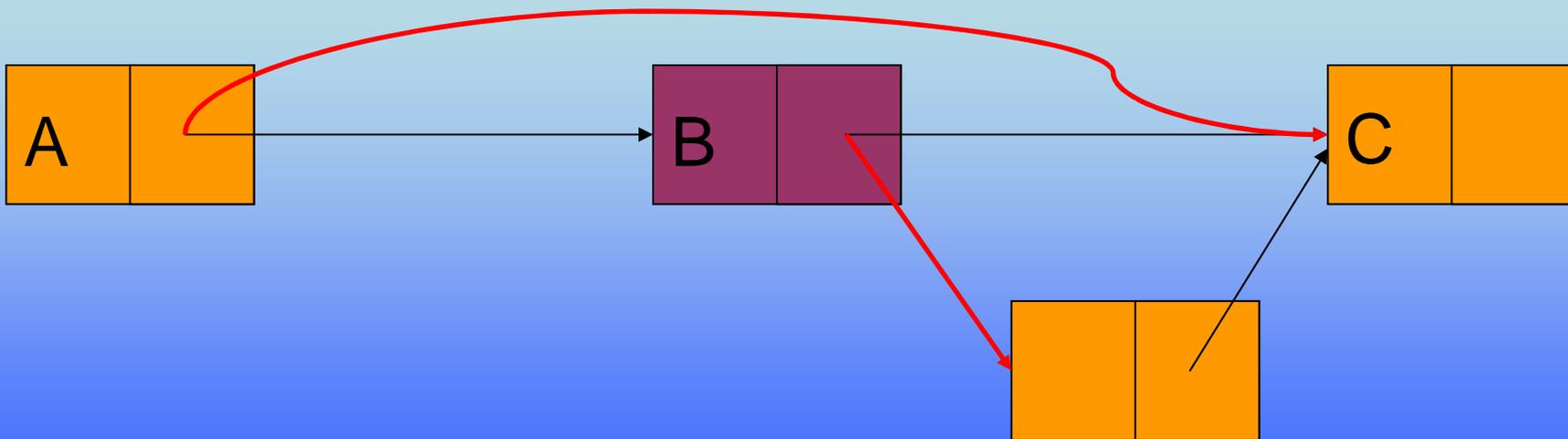
Deletion



2. The Problem

The problem

Problem: Several concurrent process => can lead to loss of data



2. The Problem

The solution: CAS

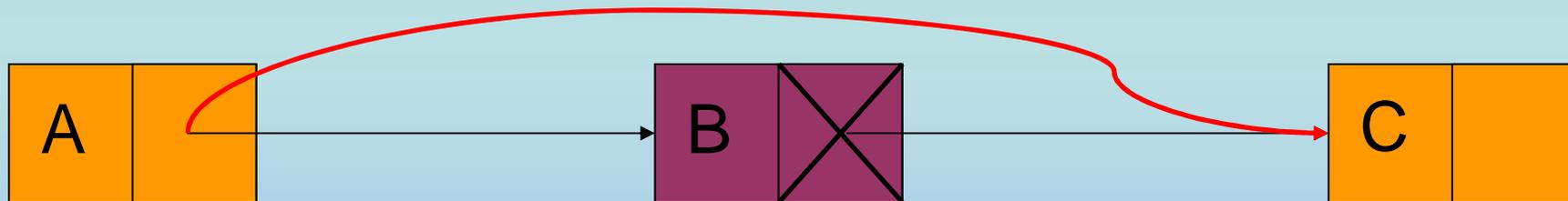
- ⇒ We do **NOT** want to change a pointer if another process change it between the time we read it and the time we write it
- ⇒ To avoid that: we need a way to verify that the pointer did not change and to change it value atomically.
- ⇒ The **Compare And Swap** atomic instruction is designed for this purpose :

```
word CAS ( word* Address, word expectedValue, word  
newValue){  
    value = *Address;  
    if ( value == expectedValue ) *Address = newValue;  
    return value;  
}
```

3. The Algorithm

The idea (1)

First solution (Harris):



A marked node can be deleted safely, as any of the process can then insert a node between B and C or delete C.

✗ Mark

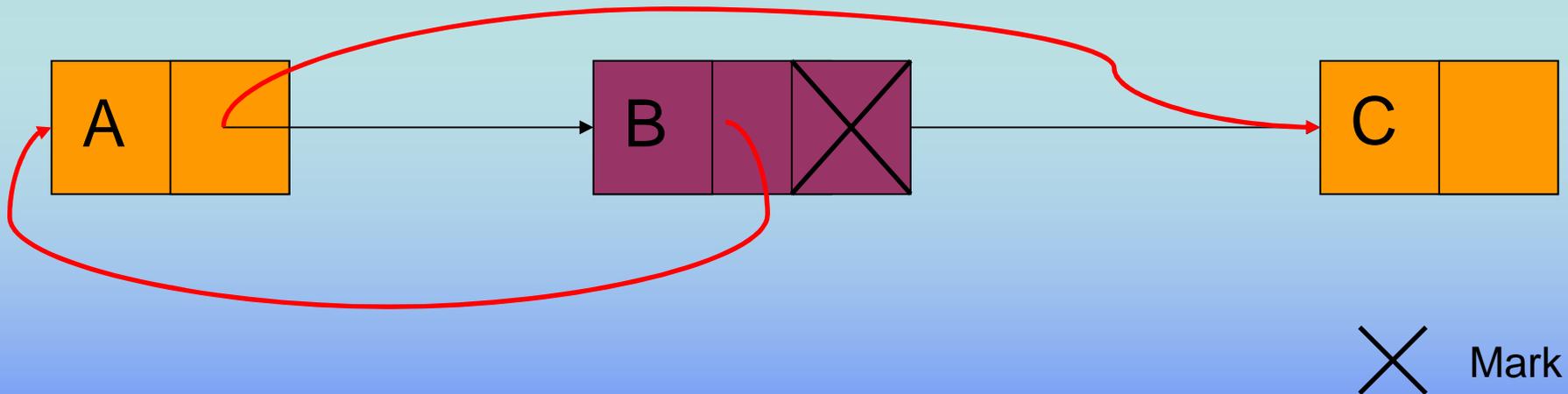
Problem:

If another process wants to delete C or insert an element between B and C, it has to re-search the element from the head !!

3. The Algorithm

The idea (2)

Improvement 1:



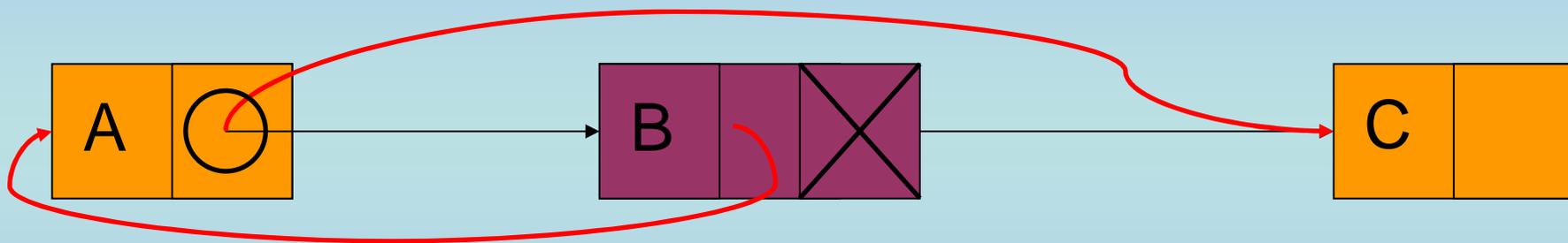
Problem:

Long back-links chain can appear !!

3. The Algorithm

The idea (3)

Improvement 2:



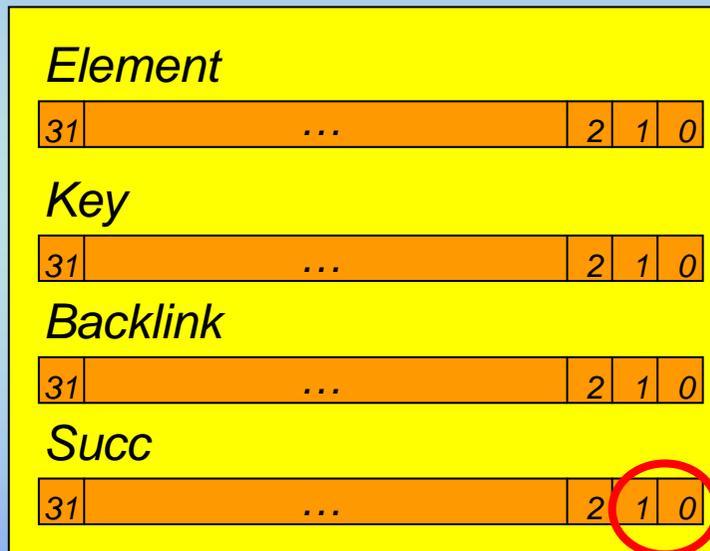
A flag prevent any operation on the node (except the deletion of the next node)



3. The Algorithm

The Data Structure

Node



Head key = $-\infty$

Tail's key = $+\infty$

2 LSB of a pointer : always 0 !

- one will represent the flag
- the other one the mark

Allow to update the 3 information atomically

3. The Algorithm

The functions

Insert

Delete

Search : just use SearchFrom

SearchFrom : used by Delete, Insert, TryFlag and Search

TryFlag : set the flag of a node

HelpFlagged : logically delete a node

TryMark : set the mark of a node

HelpMarked : physically delete a node

Try-er: loop that run until a CAS operation succeed
(additional check before CAS, update after the CAS)

Helpers: allow other process to help another one deleting
his node (pre-emptive system)

Conclusion

Algorithm complexity : n for each operation (+ contention)

Allow several process (running on different processor or the same one) to make safe operations on shared linked lists in the same time

No System Call, a process is never blocked !

No locks ! => No dead-locks !!

Useful for SMP OS design but also for any other concurrent programs

Constitute a good start point for concurrent algorithm for other data-structures

Any questions ??

Thank you for your attention!