## Chapter 1 Communication Networks and Services

### Computer Networks & Packet Switching

Alberto Leon-Garcia

Indra Widiaia

### **Computer Network Evolution Overview**



1950s - 1960s: Terminal-Oriented Computer Networks

1960s - 1970s:Computer-to-Computer Networks:the ARPANET - first Wide Area Network (WAN)

1980s:

Local Area Networks (LANs)

1980s:



### **Terminal-Oriented Networks**

- Early computer systems very expensive
- Time-sharing methods allowed multiple terminals to share local computer
- Remote access via telephone modems



### **Terminal-Oriented Networks**

#### **Example** [modulation / demodulation ]



### **Medium Access Control**



- Dedicated communication lines were expensive
- Terminals generated messages sporadically
- Frames carried messages to/from attached terminals
- Address in frame header identified terminal
- Medium Access Controls for sharing a line were developed
- Example: Polling protocol on a multidrop line



Host computer

Terminals at different locations in a city Must avoid collisions on inbound line

### **Statistical Multiplexing**



- Statistical multiplexer allows a line to carry *frames* that contain messages to/from multiple terminals
- Frames are buffered at *multiplexer* until line becomes available, i.e. store-and-forward
- Address in frame header identifies terminal
- Header carries other *control* information



### **Error Control Protocol**



- Communication lines introduced errors
- Error checking codes used on frames
  - "Cyclic Redundancy Check" (CRC) "check bits"
  - (1) CRC is calculated based on frame header and payload
  - (2) CRC is appended to frame
  - (3) if receiver detects error, retransmission is requested

CRC	Information	Header	
•			Terminal
Header	Information	CRC	



#### 1950s - 1960s:Terminal-Oriented Computer Networks

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1980s:Local Area Networks (LANs)

1980s:The Internet

### **Computer-to-Computer Networks**



- As cost of computing dropped, terminal-oriented networks viewed as too inflexible and costly
- Need to develop flexible computer networks
  - Interconnect computers as required
  - Support many applications
- Application Examples
  - File transfer between arbitrary computers
  - Execution of a program on another computer
  - Multiprocess operation over multiple computers

### **Packet Switching**



- Network should support multiple applications
  - Transfer arbitrary message size
  - Low delay for interactive applications
  - But in store-and-forward operation, long messages induce high delay on interactive messages
- Packet switching introduced
  - Network transfers packets using store-and-forward
  - Packets have maximum length
  - Break long messages into multiple packets
- ARPANET testbed led to many innovations

### **ARPANET Packet Switching**

Host generates message

Source packet switch converts message to packet(s) Packets transferred independently across network Destination packet switch reasembles message Destination packet switch delivers message



### **ARPANET Routing**

Routing is highly nontrivial in mesh networks





### **Other ARPANET Protocols**

Error control between adjacent packet switches



### **ARPANET Applications**



- ARPANET introduced many new applications
- Email, remote login, file transfer, ...
- Intelligence at the edge





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### **Local Area Networks**

#### LAN History

- in 1980s affordable computers became available
- subsequently, need for <u>low-cost</u>, high-speed, and low error-rate networks arose
  - to interconnect local workstations over small radius < 1km</li>
  - to enable sharing of local resources (printers, servers, etc.)
- complex packet switching, congestion and flow control were unnecessary
- variety of LAN topologies emerged, including: bus, ring





### Local Area Networks (cont.)

#### **Bus Topology (Ethernet)**

one long cable, so-called backbone, links all devices in the network

- each workstation connects to backbone through Network Interface Card (NIC); each NIC has globally unique address
- data frames are broadcast into coaxial cable
- receive: NIC listens to medium for frames with its address
- send: NIC listens to medium for presence of ongoing transmission if no transmission is found, send frame
- collision: if frame collides with somebody else's frame, abort transmission and retry later



### Local Area Networks (cont.)

- **Bus Topology (Ethernet)**
- advantages: simple & inexpensive installation
- **disadvantages**: 1) backbone = single point of failure
  - 2) <u>collisions  $\Rightarrow$  diminishing capacity</u>
- if two or more devices transmit simultaneously their signals will interfere





### Local Area Networks (cont.)

Ring Topology – each device has a dedicated point-to-point connection

only with the two devices on either side of it

- a small frame token circulates around the ring; only the station that possesses the token is allowed to transmit at any given time
- signal is passed along the ring in one direction, from device to device, until it reaches its destination
- advantages: fairness in access / effective use of bandwidth— token-passing provides each station with a turn to transmit
- **disadvantages**: entire network will fail if there is a failure in any transmission link or in the mechanism that relays the token







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### **The Internet**

Internet = Internetwork – two or more interconnected networks –

network of networks

#### The Internet: Past

- LANs that emerged in 1970s were different in terms of their underlying technology and operation
- a protocol that would enable communication across multiple dissimilar networks was needed
  - "higher level of abstraction" protocol
- Internet Protocol / Addressing were soon developed and enabled creation of a single global internetwork

#### **The Internet: Present**

- spread over 200 countries
- made up of 100,000s of interconnected networks, 10,000,000s of interconnected hosts, and 100,000,000s of users
- still grows exponentially ...





### The Internet (cont.)

#### **IP Network = the Internet**



- each component network must contain special packet switch, gateway / router, through which it interconnects with rest of the Internet
- host computers place data in IP packets (data + IP header) and deliver them to nearest router
- router, with help of other routers, attempts to forward packet across the Internet
- "best effort service" IP provides no mechanism to deal with packet loss, corruption, reordering



### **Addressing & Routing**

- Hierarchical address: Net ID + Host ID
- IP packets routed according to Net ID
- Routers compute routing tables using distributed algorithm





### Names and IP Addresses



- Routing is done based on 32-bit IP addresses
- Dotted-decimal notation
  - 128.100.11.1
- Hosts are also identified by name
  - Easier to remember
  - Hierarchical name structure
  - cse.yorku.ca
- Domain Name System (DNS) provided conversion between names and addresses



#### Disadvantages

- circuit establishment delay circuit establishment introduces 'initial delay'
- inefficient use of capacity channel capacity is dedicated for the duration of a connection, even if no data is being transferred (e.g. silent periods in speech)
- network complexity end-to-end circuit establishment and bandwidth allocation requires complex signaling software to coordinate operation of switches

### Packet vs. Circuit Switching (cont.)

#### Packet-Switched Networks (the Internet)



#### Advantages

greater line efficiency – network links are dynamically shared

by many packets / connections

- no blocked traffic packets are accepted even under heavy traffic. but delivery delay may increase
- Disadvantages • variable delay – each node introduces additional variable delay due to processing and queuing
  - overhead to route packets through a packet-switching network, overhead information including the address of destination and/or sequence information must be added to each packet 26