

Parallel Algorithm for Minimum Spanning Tree



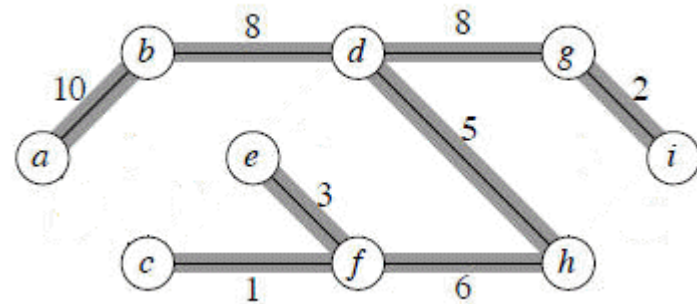
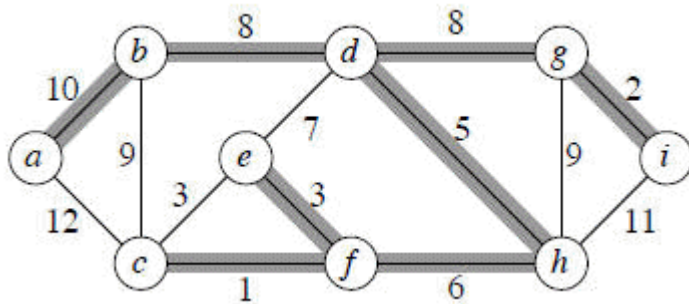
- Wrote by R. Setia, A. Nedunchezian, S. Balachandaran, in HiPC 2009
- Presented by Xiwen Chen.

Outline

- Brief introduction.
- Sequential algorithm.
- Concurrent Prim's algorithm.
- Some related problem.

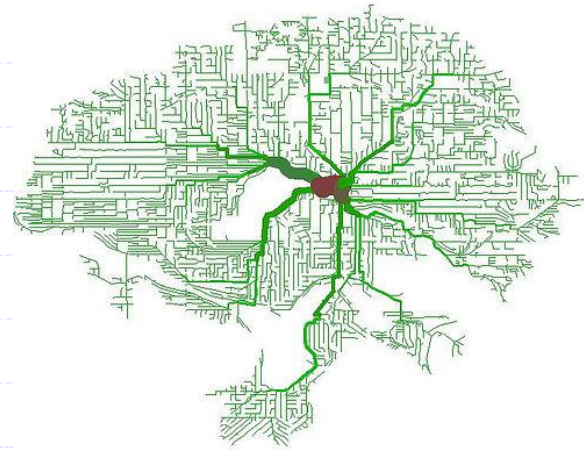
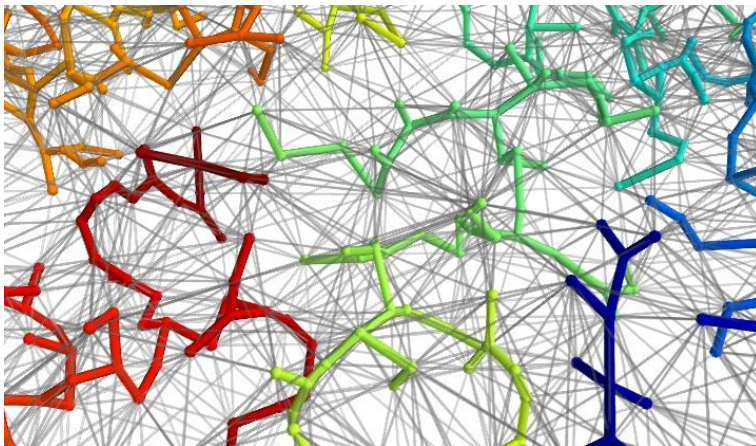
Introduction

- Minimum Spanning Tree(MST) is one of the well known classical graph problems.



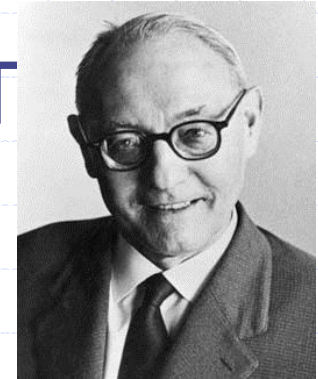
Applications

- Minimum Spanning Tree has many applications in VLSI layout and routing, wireless communication and various other fields.



Sequential Algorithms

- The first serial algorithm for MST was given by a Czech scientist Borůvka, 1926.

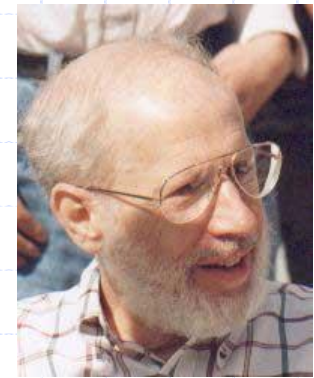


Borůvka(1899-1995)



Prim(1921-now)

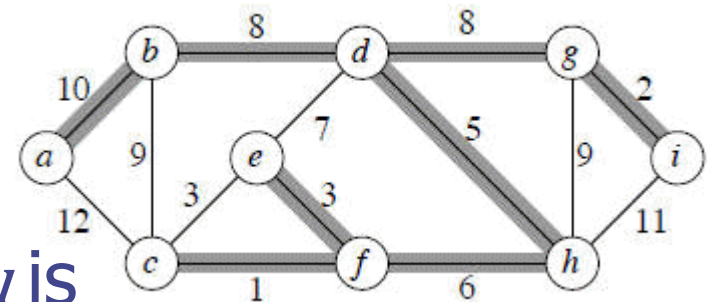
- Other two commonly used Greedy algorithms are Kruskal's algorithm (1956) and Prim's algorithm(1957).



Kruskal(1928-2010)

Prim's Algorithm

- Initialize: $V_{mst} = \{x\}$, x is an arbitrary node from V , $E_{mst} = \{\emptyset\}$.
- Repeat until $V_{mst} == V$
 - Choose an edge (u,v) with minimal weight such that u is in V_{mst} and v is not.
 - Add v to V_{mst} and (u,v) to the E_{mst} .

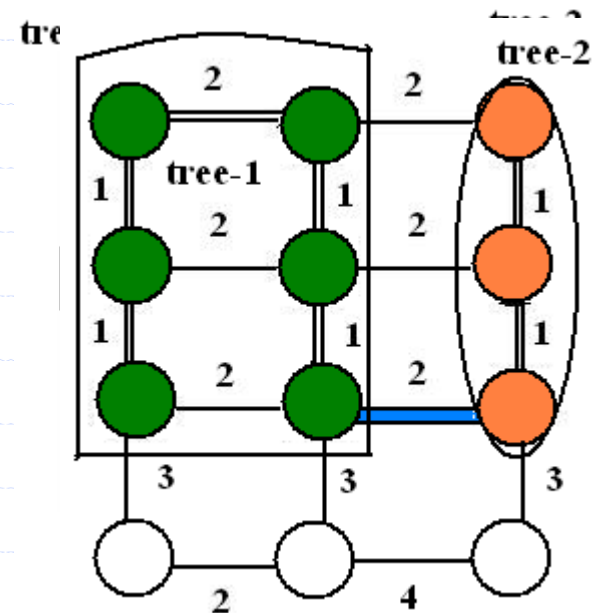


Concurrent Prim's Algorithm

- Key ideas:
 - Each thread grows/colors their own single trees in parallel.
 - When collision occurs between two threads i and j ($i < j$), thread j merges with i . i keeps growing and j chooses another node to grow a new tree.

Concurrent Prim's Algorithm

- Things remains the same for growing a single tree in an individual thread.
- For the merge part, we call the `MergeTree(i,j)` method to do that.



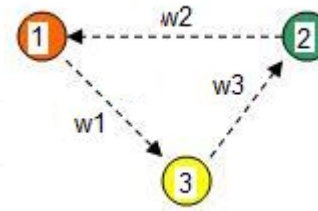
Concurrent Prim's Algorithm

- Lemma 1.1 No cycles are formed during MergeTree operation.

$w_1 < w_2$;

$w_2 < w_3$; $w_3 < w_1$; $\rightarrow w_2 < w_1$

Contradiction!



- Lemma 1.2 The edges added by the this this algorithm belong to MST.

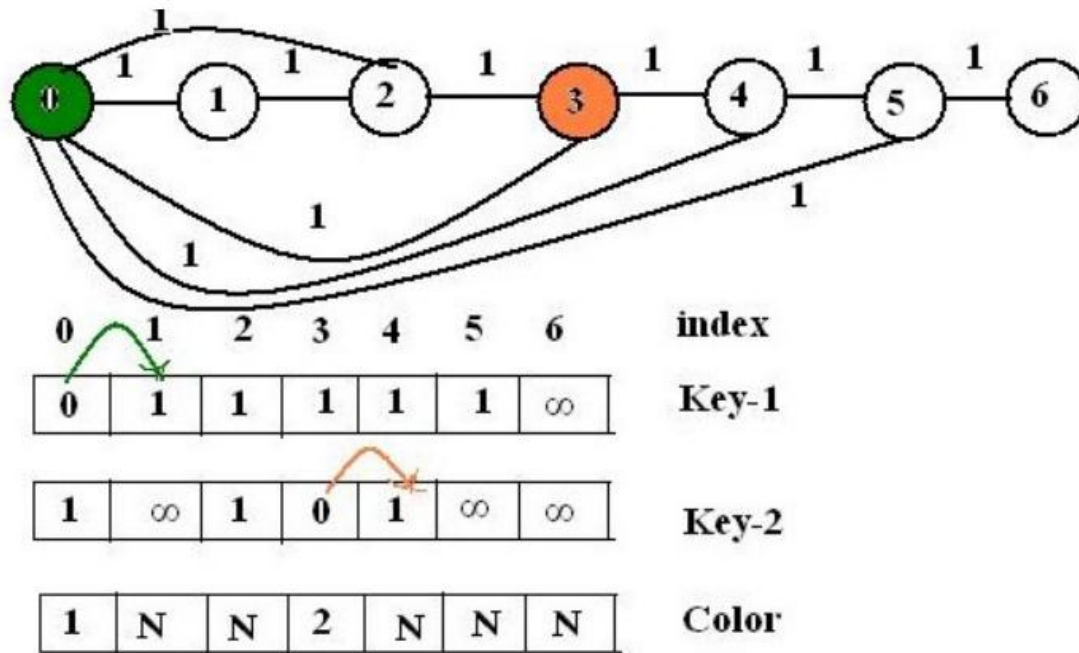
Load Balancing scheme

- Base Problem Size
 - A threshold value for the number of uncolored nodes. If # of nodes fall below that threshold, we terminate the thread j instead of let it pick a new random node to grow a new tree.

The writer has an empirical value equal to $\frac{|V|}{p}$, if there are p threads created.

Heuristics-1

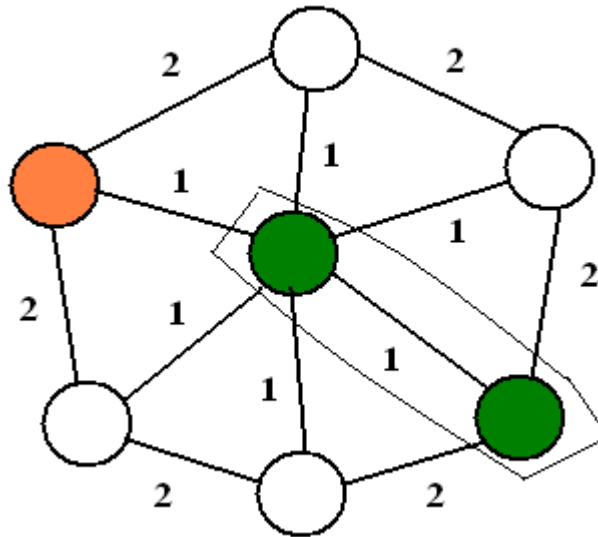
- Warp-around find-min



Heuristics-2

- Threshold nodes
 - Kill the underperformed thread

If the size of tree generated by some thread consistently smaller than N_{th} for k times, we killed that thread.



Experiments and results

Execution time vs number of processors for a given Graph and different base problem sizes

