How do you cite a journal article?

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- The names of the authors (for example, all separated by commas apart from the last two which are separated by "and").
- The title of the article.
- *The title of the journal*, volume(number): pages, month, year.

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Example

Carla Schlatter Ellis. Concurrent Search and Insertion in AVL Trees. *IEEE Transactions on Computers*, 29(9):811–817, September 1980.

How do you record a journal article in BiBTeX?

```
@article{Ellis80,
  author = "Carla Schlatter Ellis",
  title = "Concurrent Search and Insertion in
            {AVL} Trees",
  journal = "IEEE Transactions on Computers",
  volume = "29",
  number = "9",
  pages = "811--817",
  month = sep,
  year = "1980"}
```

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How do you cite a paper in a conference proceedings?

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How do you cite a paper in a conference proceedings?

- The names of the authors (for example, all separated by commas apart from the last two which are separated by "and").
- The title of the paper.
- In, the names of the editors of the proceedings (for example, all separated by commas apart from the last two which are separated by "and"), the title of the proceedings, volume of title of the series, pages, location where the conference was held, month, year.
- Publisher.

Example

Gordon Plotkin and John Power. Notions of computation determine monads. In, Mogens Nielsen and Uwe Engberg, editors, *Proceedings of the 5th International Conference on Foundations of Software Science and Computation Structures*, volume 2303 of *Lecture Notes in Computer Science*, pages 342–356, Grenoble, France, April 2002. Springer-Verlag.

Bibliography

Question

How do you record a paper in a conference proceedings in BiBTeX?

@inproceedings{PlotkinPower02,

author	= "Gordon Plotkin and John Power",
title	= "Notions of computation determine :
editor	= "Mogens Nielsen and Uwe Engberg",
booktitle	= "Proceedings of the 5th Internatio
volume	= "2303",
series	-
	= "Lecture Notes in Computer Science
pages	= "342356", "Common black Frances"
address	= "Grenoble, France",
month	= apr,
year	= "2002",
publisher	= "Springer-Verlag"}

The producer-consumer problem (also known as the bounded-buffer problem) is a classical concurrency problem.

The problem is to synchronize two threads, the producer and the consumer, who share a common, fixed-size buffer. The producer repeatedly generates a data item and puts it into the buffer. At the same time, the consumer removes data items from the buffer, one item at a time.

The problem is to make sure that the producer will not try to add data items to a full buffer and that the consumer will not try to remove data items from an empty buffer. We assume that the items are integers. We represent the buffer by means of an array of integers. The array has a fixed size.

int N = 10; // capacity of buffer

The producer and consumer share the following variables.

```
int[] buffer; // array representing buffer
int next = 0; // index of cell for next item
int size = 0; // number of items stored in buffer
```

Producer:

```
while (true)
int value = produce an item;
buffer[next] = value;
size++;
next = (next + 1) mod N;
```

Consumer:

```
while (true)
    int value = buffer[(next - size) mod N];
    size--;
```

How can we make sure that the producer will not try to add data items to a full buffer?

Question

How can we make sure that the consumer will not try to remove data items from an empty buffer?

The readers and writers problem, due to Courtois, Heymans and Parnas, is another classical concurrency problem. It models access to a database. There are many competing threads wishing to read from and write to the database. It is acceptable to have multiple threads reading at the same time, but if one thread is writing then no other thread may either read or write. The problem is how do you program the reader and writer threads? The readers and writers share the following variable. semaphore mutex = 1; Reader: P(mutex); read; V(mutex); Writer: P(mutex); write; V(mutex);

Does it solve the readers-writers problem?

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Does it solve the readers-writers problem?

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Answer

Yes!

Does it solve the readers-writers problem?

Answer

Yes!

Question

Is it a satisfactory solution?

Does it solve the readers-writers problem?

Answer

Yes!

Question

Is it a satisfactory solution?

Answer

No!

The Readers-Writers Problem

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Question

Why not?

Why not?

Answer

It does not allow multiple readers to read at the same time.

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While a writer is writing, readers and writers arrive. Once the writer is done, can the readers start reading?

Options

While readers are reading, readers and writers arrive. Can the readers start reading?

While a writer is writing, readers and writers arrive. Once the writer is done, can the readers start reading?

Yes

Options

While readers are reading, readers and writers arrive. Can the readers start reading?

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Yes

No reader is kept waiting unless a writer is writing.

While a writer is writing, readers and writers arrive. Once the writer is done, can the readers start reading?

Options

While readers are reading, readers and writers arrive. Can the readers start reading?

While a writer is writing, readers and writers arrive. Once the writer is done, can the readers start reading?

No

Options

While readers are reading, readers and writers arrive. Can the readers start reading?

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No

If a writer wants to write, it writes as soon as possible.

In the dining philosophers problem, due to Dijkstra, five philosophers are seated around a round table. Each philosopher has a plate of spaghetti. The spaghetti is so slippery that a philosopher needs two forks to eat it. The layout of the table is as follows.



The life of a philosopher consists of alternative periods of eating and thinking. When philosophers get hungry, they try to pick up their left and right fork, one at a time, in either order. If successful in picking up both forks, the philosopher eats for a while, then puts down the forks and continues to think.

Edsger Wybe Dijkstra

- Member of the Royal Netherlands Academy of Arts and Sciences (1971)
- Distinguished Fellow of the British Computer Society (1971)
- Recipient of the Turing Award (1972)
- Foreign Honorary Member of the American Academy of Arts and Sciences (1975)



Edsger Wybe Dijkstra

(1930-2002)

```
int N = 5;
philosopher(i):
while (true)
  think;
  takeForks(i);
  eat;
  putForks(i);
takeForks(i):
putForks(i):
```

Question

Is this solution correct?



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Question

Is this solution correct?

Answer

No!

Question
Is this solution correct?
Answer
No!
Question
Why not?



Question	
Is this solution correct?	
Answer	
No!	
Question	
Why not?	

Answer

Because philosopher(0) and philosopher(1) cat eat at the same time.

```
int N = 5;
semaphore mutex = 1;
philosopher(i):
while (true)
  think:
  takeForks(i);
  eat;
  putForks(i);
takeForks(i):
P(mutex);
putForks(i):
V(mutex);
```

Question

Is this solution correct?



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Question

Is this solution correct?

Answer

Yes!

Question

Is this solution correct?

Answer

Yes!

Question

Does this solution allow two philosophers to eat at the same time?

Question	
Is this solution correct?	
Answer	

Yes!

Question

Does this solution allow two philosophers to eat at the same time?

Answer

No!

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```
int N = 5;
int state [];
semaphore mutex = 1;
philosopher(i):
while (true)
  think;
  takeForks(i);
  eat;
  putForks(i);
```

```
takeForks(i):
P(mutex);
if (state [(i - 1) mod N] != EATING
    and state [(i + 1) mod N] != EATING)
  state[i] = EATING;
else
  state[i] = HUNGRY;
V(mutex):
putForks(i):
P(mutex);
state[i] = THINKING;
V(mutex);
```

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