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W4-M

System Calls - Examples

Example # 1: Print a string

```
.data
    str:    .asciiiz "the answer is"
.text
    addi $v0,$zero,4
    la $a0,str    # pseudoinstruction
    syscall
```

Example # 2: Input an integer

```
addi $v0,$zero,5
syscall
```

Example # 3: Print an integer

```
addi $v0,$zero,1
addi $a0,$s0,$zero
syscall
```

Example # 4: Read String

```
addi $v0,$zero,8
la $a0,Buff  #$a0=address of Buff
addi $a1,$zero,60 #$a1=max. len.
syscall
```

Service	System Call Code (\$v0)	Arguments	Result
print_int	1	\$a0 = int	
print_float	2	\$f12 = float	
print_double	3	\$f12 = double	
print_string	4	\$a0 = string address	
read_int	5		int (in \$v0)
read_float	6		float (in \$f0)
read_double	7		double (in \$f0)
read_string	8	\$a0 = buffer \$a1 = length	
sbrk	9	\$a0 = amount	address (in \$v0)
exit	10		terminate prog

Putting it all together (2)

Activity: Write a MIPS program which does the following:

1. Accepts an integer N using the following prompt
Please input a value for N =
2. Computes the sum of integers from 1 to N, i.e., $(1 + 2 + \dots + N)$ if $N > 0$
3. Displays the result (X) as
The sum of the integers from 1 to N is X
2. Waits for the next number N.
3. If $N \leq 0$, the program exits with the following farewell
Chao - Have a good day

Run the program in the spim simulator to verify the results

main routine

```
1 .data
2     str1:    .asciiiz "\nPlease input a value for N = "
3     str2:    .asciiiz "The sum of the integers from 1 to N is...\\n\\t\\t\\t"
4     str3:    .asciiiz "\\n\\\"Chao\\\" - Have a good day"
5 .text
6 main:
7     addi $v0,$zero,4
8     la $a0,str1      # pseudoinstruction
9     syscall          # Print out user prompt
10    #-----
11    addi $v0,$zero,5  # input integer
12    syscall
13    #-----
14    add $a0,$zero,$v0  # Move input integer as to $a0 - input to sum
15    slt $t0,$v0,$zero
16    bne $t0,$zero,end  # if input <0 end
17    jal sum            # Jump to subroutine sum
18    #-----
19    addi $sp, $sp, -4
20    sw $v0, 4($sp)      # Put the sum on the stack
21    addi $v0,$zero,4
22    la $a0,str2
23    syscall          # Print message for sum
24    #-----
25    addi $v0, $zero,1
26    lw $a0, 4($sp)      # Get sum from stack
27    addi $sp, $sp,4
28    syscall          # Print sum
29    #-----
30 end:
31     addi $v0,$zero,4
32     la $a0,str3
33     syscall          # Print out good bye message
34     add $a0,$a0,$zero
35     addi $v0,$zero,17
36     syscall          # exit2
37    #-----
```

sum subroutine

```
38 sum:
39     addi $sp, $sp, -8  # Set up the stack
40     sw $ra, 4($sp)      # Save return address
41     addi $t0, $a0, 0      # Initialize the sum
42     li $v0, 0            # Initialize return value
43     beq $t0,$zero, ret  # If argument is 0 then return
44     addi $t1, $t0, -1      # Compute n-1
45     sw $t0, 8($sp)      # Save caller saved regs
46     addi $a0, $t1, 0      # Move n-1 into argument register
47     jal sum              # Call sum
48    #-----
49     lw $t0, 8($sp)      # Restore caller saved reg
50     add $v0, $t0, $v0      # Add return value to $t0
51     lw $ra, 4($sp)      # Get the return address
52     addi $sp,$sp,8
53 ret:
54     jr $ra                # Return
```

Same Example – More Verbose Comments

```

1 ##########
2 # Program Name: Sum of Positive Integers
3 # Programmer: H. Chesser
4 # Date last Modified: Sept 24, 2012
5 #########

6

7

8 #########
9 # Procedure <main> #
10 #########

11

12 # The assembler directive .data tells the assembler that all data allocation directives should
13 # allocate data in a portion of memory called the data segment
14

15     .data
16 Prompt:    .asciiz "\n Please input a value for N = "
17 Result:    .asciiz "The sum of the integers from 1 to N is... \n\t\t\t"
18 Bye:       .asciiz "\n*** \'Chao\' - Have a good day *** \n"
19

20 # The assembler directive .globl <label> declare that <label> is global so it can be referenced
21 # from other files
22     .globl main
23

24 # The assembler directive .text <address> tells the assembler that subsequent items must be
25 # placed in the user text segment. In SPIM, these items may only be instructions or words. If
26 # the optional argument <address> is specified, subsequent items are stored starting at address
27 # <address>
28     .text
29 ##########
30 main:
31     addi $sp, $sp, -12
32     sw $ra, 8($sp)
33     sw $a0, 4($sp)
34     sw $v0, 0($sp)
35

36 ##########
37 again:   addi $v0, $zero, 4      # Sys Call # 4 (Print String)
38     la $a0, Prompt        # load address of Prompt into $a0
39     syscall               # print the Prompt message
40

41     addi $v0, $zero, 5      # Sys Call # 5 (Read Integer)
42     syscall               # reads the value of N into $v0
43

44     slti $t0, $v0, 1       # Set $t0 = 1 if $v0 < 1, else $t0 = 0
45     bne $t0, $zero, End    # if ($v0 < 1), go to End
46

47     add $a0, $v0,$zero      # $v0 = $a0
48     jal myadd              # call myadd
49

50     addi $t1, $v0, 0        # $t1 = $v0
51     addi $v0, $zero, 4      # Sys Call # 4 (Print String)
52     la $a0, Default        # load address of Default into $a0

```

```

    . . . .

75 ##########
76 # Procedure <myadd> #
77 #########

78

79

80 ########## save registers #####
81 myadd:
82     addi $sp, $sp, -8
83     sw $ra, 4($sp)
84     sw $s1, 0($sp)
85

86 ########## myadd body #####
87     addi $s1, $zero, 0      # $s1 = 0
88 Loop:
89     add $s1, $s1, $a0      # $s1 = $s1 + $a0
90     addi $a0, $a0, -1       # $a0 = ($a0 - 1)
91     slti $t0, $a0, 1        # $t0 = 1 if $a0 < 1, else $t0 = 0
92     beq $t0, $zero, Loop    # if ($v0 > 1), go to Loop
93     addi $v0, $s1, 0
94

95 ########## restoring registers #####
96     lw $s1, 0($sp)
97     lw $ra, 4($sp)
98     addi $sp, $sp, 8
99

100 ########## procedure return #####
101    jr $ra
102

```

Agenda for Today

1. Addition, Subtraction
2. Overflow
3. Multiplication
4. Division
5. Floating Point: IEEE 754 single and double precision formats

Patterson: Sections 3.1 – 3.5

Wednesday – Quiz #1

Addition and Subtraction

In MIPS, addition and subtraction for signed numbers use 2's complement arithmetic

Example 1: Add 10_{ten} and 15_{ten}

bit 1	bit 2	Prev. Carry	Sum	Next Carry
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

Table 1: Truth Table for addition

Example 2: Subtract 15_{ten} from 10_{ten}

Addition and Subtraction

In MIPS, addition and subtraction for signed numbers use 2's complement arithmetic

Example 1: Add 10_{ten} and 15_{ten}

Step 1: Represent the operands in 2's complement

$$10_{\text{ten}} = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1010_{\text{two}}$$

$$15_{\text{ten}} = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1111_{\text{two}}$$

Step 2: Perform bit by bit addition using table 1.

$$25_{\text{ten}} = 0000\ 0000\ 0000\ 0000\ 0000\ 0001\ 1001_{\text{two}}$$

Example 2: Subtract 15_{ten} from 10_{ten}

The problem is reduced to $(10_{\text{ten}} + (-15_{\text{ten}}))$

$$10_{\text{ten}} = 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1010_{\text{two}}$$

$$-15_{\text{ten}} = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 0001_{\text{two}}$$

$$-5_{\text{ten}} = 1111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1011_{\text{two}}$$

bit 1	bit 2	Prev. Carry	Sum	Next Carry
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

Table 1: Truth Table for addition

Overflow (1)

Recall that:

Smallest signed integer:

$$\begin{aligned} &1000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000_2 \\ &= -(2^{31})_{10} = -2,147,483,648_{10} \end{aligned}$$

Largest signed integer:

$$\begin{aligned} &0111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1111_2 \\ &= (2^{31} - 1)_{10} = 2,147,483,647_{10} \end{aligned}$$

What happens if the result of an operation is more than the largest signed integer or less than the smallest signed integer?

Example: Add $2,147,483,640_{10}$ and 28_{10}

$$\begin{array}{r} 28_{10} \\ + 2,147,483,640_{10} \\ \hline \end{array} \quad \begin{array}{l} = 0000\ 0000\ 0000\ 0000\ 0000\ 0001\ 1100_2 \\ = 0111\ 1111\ 1111\ 1111\ 1111\ 1111\ 1000_2 \end{array}$$

$$\begin{array}{r} 28_{10} + 2,147,483,640_{10} \\ = 1000\ 0000\ 0000\ 0000\ 0000\ 0001\ 0100_2 \\ = -2,147,483,628_{10} \end{array}$$

Overflow caused the value to be perceived as a negative integer

Overflow (2)

When can overflow occur?

Operation	Operand A	Operand B	Result indicating overflow
$A + B$	$A \geq 0$	$B \geq 0$	< 0
$A + B$	$A < 0$	$B < 0$	≥ 0
$A - B$	$A \geq 0$	$B < 0$	< 0
$A - B$	$A < 0$	$B \geq 0$	≥ 0

Integer Multiplication, Division

Both operations really imply a series of additions and subtractions

Example: Multiply 10_{ten} and 3_{ten} :

$$10_{\text{ten}}$$

$$= 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1010_{\text{two}}$$

$$3_{\text{ten}}$$

$$= 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0011_{\text{two}}$$

$$10_{\text{ten}} * 1_{\text{ten}}$$

$$= 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1010_{\text{two}}$$

$$10_{\text{ten}} * 2_{\text{ten}}$$

$$= \underline{0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0001\ 0100}_{\text{two}}$$

$$= 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 0001\ 1110_{\text{two}} = 0x1e = 30_{\text{ten}}$$

Example: Divide 28_{ten} by 9_{ten} :

$$28_{\text{ten}}$$

$$= 0000\ 0000\ 0000\ 0000\ 0000\ 0001\ 1100_{\text{two}}$$

$$9_{\text{ten}}$$

$$= 0000\ 0000\ 0000\ 0000\ 0000\ 0000\ 1001_{\text{two}}$$

$$\begin{array}{r} 11 \\ 1001 \overline{)11100} \\ \underline{1001} \\ 1010 \\ \underline{1001} \\ 1 \end{array}$$

Floating Point: Single Precision

1. In MIPS, decimal numbers are represented with the [IEEE 754 binary representation](#) that uses the **normalized** standard scientific binary notation defined as

$$(-1)^S \times (1 + \text{fraction})_{\text{two}} \times 2^{\text{exponent} - \text{bias}}$$

2. A number in **normalized scientific notation** has a mantissa that has no leading 0's and must be of the form $(1 + \text{fraction})$. For example, the binary representations 2.0×2^{-5} , 0.5×2^{-3} , 4.0×2^{-6} , and 1.0×2^{-4} are all equivalent but only 1.0×2^{-4} is the normalized scientific binary notation.
 3. MIPS allows for two floating point representations: Single precision and double precision.
 4. **Single precision** has a bias of 127 while double precision has a bias of 1023.
 5. In single precision, the floating point representation is 32 bit long and has the following form

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

S exponent

fraction

(8 bits)

(23 bits)

where S represents the sign bit, which is 1 for negative numbers and 0 for positive numbers.

Activity 2:

Represent -0.75_{ten} , 1.4_{ten} in single precision of IEEE 754 binary representation.

Floating Point: Double Precision

1. In **double precision**, the value of bias in
$$(-1)^S \times (1 + \text{fraction})_{\text{two}} \times 2^{\text{exponent} - \text{bias}}$$
is 1023.
 2. In single precision, the floating point representation is 32 bit long and has the following form

31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0

S	exponent (11 bits)	fraction (Total of 52 bits)
----------	-------------------------------------	--

fraction (continued)

Activity 3:

Represent -0.75_{ten} in double precision of IEEE 754 binary representation.

Activity 4:

Show that the largest magnitude that can be represented using single precision is $\pm 6.8_{\text{ten}} \times 10^{38}$, while the smallest fraction that can be represented is $\pm 5.9_{\text{ten}} \times 10^{-39}$.

Floating Point Registers

Name	Example	Comments
32 floating point registers each is 32 bits long	\$f0, \$f1, \$f2, \$f3, \$f4 ,..., \$f31	MIPS floating point registers are used in pairs for double precision numbers
Memory w/ 2^{30} words	Memory[0], Memory[4], ... Memory[4294967292]	Memory is accessed one floating point (single or double precision) at a time

The following is the established register usage convention for the floating point registers:

- \$f0, \$f1, \$f2, \$f3:** Function-returned values
- \$f4, \$f5, ..., \$f11:** Temporary values
- \$f12, \$f13, \$f14, \$f15:** Arguments passed into a function
- \$f16, \$f17, \$f18, \$f19:** More Temporary values
- \$f20, \$f21, ..., \$f31:** Saved values

Floating Point Instructions

Category	Instruction	Example	Meaning	Comments
Arithmetic	FP add single	add.s \$f2,\$f4,\$f6	$\$f2 \leftarrow \$f4 + \$f6$	Single Prec.
	FP subtract single	sub.s \$f2,\$f4,\$f6	$\$f2 \leftarrow \$f4 - \$f6$	Single Prec.
	FP multiply single	mul.s \$f2,\$f4,\$f6	$\$f2 \leftarrow \$f4 \times \$f6$	Single Prec.
	FP divide single	div.s \$f2,\$f4,\$f6	$\$f2 \leftarrow \$f4 / \$f6$	Single Prec.
	FP add double	add.d \$f2,\$f4,\$f6	$\$f2 \leftarrow \$f4 + \$f6$	Double Prec.
	FP subtract double	sub.d \$f2,\$f4,\$f6	$\$f2 \leftarrow \$f4 - \$f6$	Double Prec.
	FP multiply double	mul.d \$f2,\$f4,\$f6	$\$f2 \leftarrow \$f4 \times \$f6$	Double Prec.
	FP divide double	div.d \$f2,\$f4,\$f6	$\$f2 \leftarrow \$f4 / \$f6$	Double Prec.
Data Transfer	load word copr.1	lwc1 \$f2,100(\$s2)	$\$f2 \leftarrow \text{Mem}[\$s2+100]$	Single Prec.
	store word copr.1	swc1 \$f2,100(\$s2)	$\text{Mem}[\$s2+100] \leftarrow \$f2$	Single Prec.
	FP compare single (eq, ne, lt, le, gt, ge)	c.lt.s \$f2,\$f4	if ($\$f2 < \$f4$) cond = 1, else cond = 0	Single Prec.
	FP compare double (eq, ne, lt, le, gt, ge)	c.lt.d \$f2,\$f4	if ($\$f2 < \$f4$) cond = 1, else cond = 0	Double Prec.
Conditional branch	Branch on FP true	bclt 25	if cond==1 go to PC+100+4	Single/ Double Prec.
	Branch on FP false	bclf 25	if cond==0 go to PC+100+4	Single/ Double Prec.

Example

```
# calculate area of a circle

        .data

Ans:    .asciiiz      "The area of the circle is: "
Ans_add:.word      Ans                      # Pointer to String (Ans)
Pi:     .double       3.1415926535897924
Rad:    .double       12.345678901234567
Rad_add:.word      Rad                      # Pointer to float (Rad)

        .text

main:   lw $a0, Ans_add($0)          # load address of Ans into $a0
        addi $v0, $0, 4                # Sys Call 4 (Print String)
        syscall

#-----                         # load float (Assembler Instruction)
        la $s0, Pi                  # load address of Pi into $s0
        ldc1 $f2, 0($s0)            # $f2 = Pi

#-----                         # load float (MIPS Instruction)
        lw $s0, Rad_add($0)          # load address of Rad into $s0
        ldc1 $f4, 0($s0)            # $f4 = Rad
        mul.d $f12, $f4, $f4
        mul.d $f12, $f12, $f2
        addi $v0, $0, 3              # Sys Call 3 (Print Double)
        syscall

exit:  jr $ra
```