CSE 3214: Computer Network Protocols and Applications -Transport Layer

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Chapter 3: Transport Layer

our goals:

- understand principles behind transport layer services:
 - multiplexing, demultiplexing
 - reliable data transfer
 - flow control
 - congestion control

- learn about Internet transport layer protocols:
 - UDP: connectionless transport
 - TCP: connection-oriented reliable transport
 - TCP congestion control

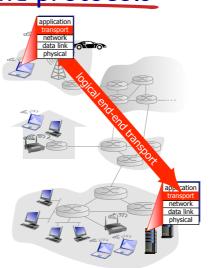
Chapter 3 outline

- 3.1 transport-layer services
- 3.2 multiplexing and demultiplexing
- 3.3 connectionless transport: UDP
- 3.4 principles of reliable data transfer
- 3.5 connection-oriented transport: TCP
 - segment structure
 - reliable data transfer
 - flow control
 - connection management
- 3.6 principles of congestion control
- 3.7 TCP congestion control

Transport Layer 3-3

Transport services and protocols

- provide logical communication between app processes running on different hosts
- transport protocols run in end systems
 - send side: breaks app messages into segments, passes to network layer
 - receive side: reassembles segments into messages, passes to app layer
- more than one transport protocol available to apps
 - Internet: TCP and UDP



Transport vs. network layer

- network layer: logical communication between hosts
- transport layer: logical communication between processes
 - relies on, enhances, network layer services

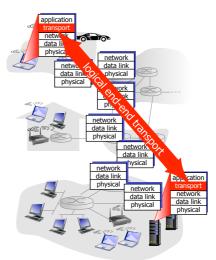
household analogy:

- 12 kids in Ann's house sending letters to 12 kids in Bill's house:
- hosts = houses
- processes = kids
- app messages = letters in envelopes
- transport protocol = Ann and Bill who demux to inhouse siblings
- network-layer protocol = postal service

Transport Layer 3-5

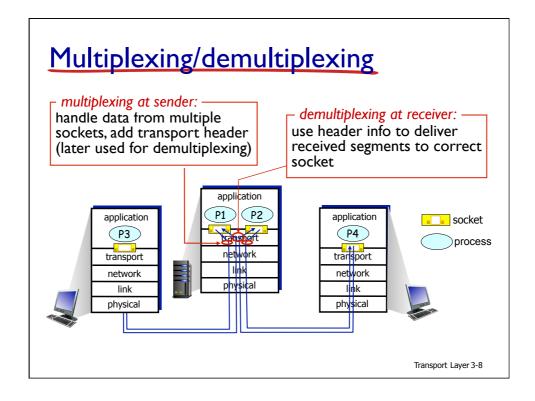
Internet transport-layer protocols

- reliable, in-order delivery (TCP)
 - congestion control
 - flow control
 - connection setup
- unreliable, unordered delivery: UDP
 - no-frills extension of "best-effort" IP
- services not available:
 - delay guarantees
 - bandwidth guarantees



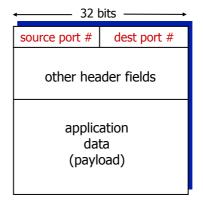
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How demultiplexing works

- host receives IP datagrams
 - each datagram has source IP address, destination IP address
 - each datagram carries one transport-layer segment
 - each segment has source, destination port number
- host uses IP addresses & port numbers to direct segment to appropriate socket



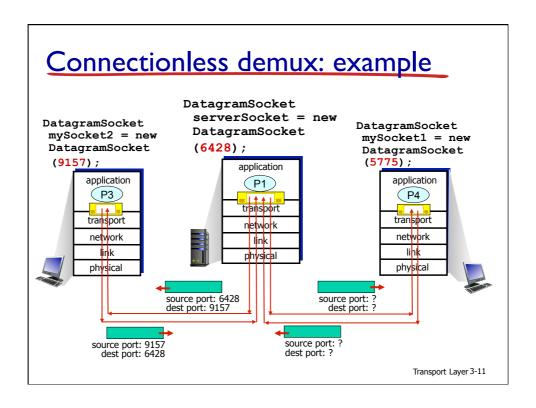
TCP/UDP segment format

Transport Layer 3-9

Connectionless demultiplexing

- recall in Socket Programming: clientSocket=socket(AF_INET, SOCK_DGRAM)
- clientSocket.sendto(message,(serverName, serverPort))
- When creating datagram to send into UDP socket, must specify
 - destination IP address
 - destination port #
- when host receives UDP segment:
 - checks destination port # in segment
 - directs UDP segment to socket with that port #

IP datagrams with same dest IP addr, & dest. port #, but different source IP addresses and/or source port numbers will be directed to same socket at dest



Connection-oriented demux

- TCP socket identified by 4-tuple:
 - source IP address
 - source port number
 - dest IP address
 - dest port number
- demux: receiver uses all four values to direct segment to appropriate socket
- server host may support many simultaneous TCP sockets:
 - each socket identified by its own 4-tuple
- web servers have different sockets for each connecting client
 - non-persistent HTTP will have different socket for each request

