

Prolog Core Concepts and Notation

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Adapted from Peter Roosen-Runge

Readings: C & M Ch 1, 2, 3.1-3.3, 8

declarative/logic programming

- idea: write a program that is a logical theory about some domain and then query it
- most well known instance is Prolog
- core constructs, terms and statements, are inherited from first order logic

terms

- Prolog statements express relationships among *terms*
- terms are (a generalization) of the same notion in first order logic, i.e. a constant, a variable, or a function applied to some argument terms
- E.g. `john`, `john_smith`, `X`, `Node`, `_person`, `fatherOf(paul)`, `date(25,10,2005)`
- `fatherOf` and `date` are functors; `date` has arity 3; it takes 3 arguments

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terms

- variables begin with upper-case letter or `_`
- constants and functors (symbols) begin with lower-case
- terms denote objects
- compound terms are called structures
- E.g. `course(complexity,time(Monday,9,11),lecturer(patrick,dymond),location(CSE,3311))`
- used to represent complex data
- terms (usually) have a tree structure

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facts

- *facts* are like atomic formulas in first order logic.
- syntax is same as terms, but ending with a period.
- e.g. `fatherOf(paul,henry).`
`mortal(ulyssus).`
`likes(X,iceCream).`
`likes(mary,brotherOf(helen)).`
- variables are implicitly universally quantified.

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rules

- *rules* are conditional statements.
- e.g. `mortal(X) :- human(X).`
i.e. $\forall x \text{ Human}(x) \rightarrow \text{Mortal}(x)$,
all humans are mortal.
- `daughter(X,Y) :- father(Y,X), female(X).`
- `,` represents conjunction.
- `likes(mary,X) :- isSweet(X).`

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rules

- `ancestor(X,Y) :- father(X,Z), ancestor(Z,Y).`
- variables are universally quantified from outside; can think of variables that appear only in rule body as existentially quantified.

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queries

- A query asks whether a given statement is true, i.e. whether it follows from the program.
- e.g. `?- mortal(ulyssus).` given
`mortal(X) :- human(X).`
`human(ulyssus).` `human(penelope).`
`god(zeus).`
Prolog answers **Yes**

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queries

- `?- mortal(X).`
`X = ulyssus ;`
`X = penelope`
`Yes`
- variables in queries are existentially quantified; can be used to retrieve information.
- can have conjunctive queries, e.g.
`?- mortal(X), mortal(Y), not(X = Y).`

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lists

- lists are a special kind of term that allows arbitrary number of components
- `[]` is the empty list
- `.(a,b)` is a dotted pair
- `[a, b, c] = .(a,.(b,.(c,[])))` is a list of 3 components.
- the functor `.` builds binary trees
- can use `display(X)` to print internal representation of X

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lists

- can refer to the first and rest of a list using the notation: `[First | Rest]`
- e.g. `?- X = [a,b,c], X = [F|R].`
`X = [a,b,c]`
`F = a`
`R = [b,c]`
- E.g. `X = [b], Y = a, Z = [Y|X].`
`X = [b]`
`Y = a`
`Z = [a,b]`

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unification

- this was an instance of the kind of pattern matching called **unification** that Prolog performs
- Prolog tries to find a way to instantiate the variables (substitute terms for them) that satisfies the query
- more on this later

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terms can represent graphs

- ?- $X = [a|X]$.
 $X = [a, a, a, a, a, a, a, a, a|...]$
Yes
- here X denotes an infinite or circular list
- this is not allowed in first-order logic; a variable cannot denote a term and one of its subterms; but Prolog omits the “occurs check”

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building a knowledge base

- to be used in a computation, facts and rules must be stored in the (dynamic) database
- facts and rules get into the database through *assertion* and *consultation*
- consultation loads facts and rules from a file

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assertion

- `?- assert(human(ulyssus)).`
- `?- human(X).`
X = ulyssus
Yes
- assertion can be done dynamically
- also `retract` to remove facts and rules from the DB
- like assignment, change state; avoid when possible

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consultation

- `?- consult('family.pl').`
loads facts and rules from file family.pl
- `?- [family].`
does the same thing
- `?- [user].`
lets you enter facts and rules from the keyboard

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denotation/meaning of Prolog programs

- a Prolog program defines a set of relations, i.e. specifies which tuples of objects/terms belong to a particular relation
- in logic, this is called a model
- declarative programming is very different from usual procedural programming where programs perform many state changing operations

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denotation of Prolog program e.g.

- `fatherOf(john,paul).`
`fatherOf(mary,paul).`
`motherOf(john,lisa).`
`parentOf(X,Y) :- fatherOf(X,Y).`
`parentOf(X,Y) :- motherOf(X,Y).`
- `fatherOf` is the relation {<john,paul>, <mary,paul>}
- what is the relation associated with `motherOf` and `parentOf`?

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rules as procedures

- rule has form goal :- body
- goal or head is like name of procedure
- terms on the RHS are like the body of the procedure, the sub-goals that have to be achieved to show that the goal holds
- the sub-goals will be attempted left-to-right
- rule succeeds if all sub-goals succeed

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passing values

- calling/querying a goal can instantiate its variables
- a sub-goal's success can bind a variable within it, also binding the same variable in the goal
- binding or instantiating a variable is giving it a value
- compare to passing values into or out of a procedure

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passing values e.g.

- Assume program:
motherOf(john,lisa).
parentOf(X,Y) :- motherOf(X,Y).
- Queries:
?- parentOf(john,X).
X = lisa Yes
?- parentOf(X,lisa).
X = john Yes
?- parentOf(X,Y).
X = john, Y = lisa Yes
- No fixed input and output parameters

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almost everything is syntactically a term

- lists are terms; what is the functor?
- rules are terms:
grandfather(X,Y):- father(X,Z),
father(Z,Y).
What are the functors?
- queries are terms

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operators

- some functors are represented by *infix* or *prefix* or *postfix* operators
- Some infix operators: `is`, `=`, `+`, `*`, `/`, `mod`, `>`, `>=`, `:-`, `,`, etc.
- `+` and `-` are both prefix and infix
- `:-` as prefix is a command, used for declarations
- operators have precedence
- can define our own operators

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arithmetic functions

- Prolog retains arithmetic functions as functions (more intuitive):
`?- X is exp(1). % exp(1) = e1`
`X = 2.71828`
`Yes`
`?- X is (4 + 2) * 5.`
`X = 30`
`Yes`
- How does `is` compare with `=`, assignment?

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relational thinking

- in Prolog, formulate statements about function values as relational facts, e.g. `factorial(0,1)`.
`factorial(N,M):- K is N -1, factorial(K,L),
M is N * L.`
- to compose functions, e.g. $Y = f(g(X))$, you must name intermediate results
`fg(X,Y):- g(X,Z), f(Z,Y)`.

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help is sometimes helpful

?- `help(reverse)`.

`reverse(+List1, -List2)`

Reverse the order of the elements in List1 and unify the result with the elements of List2.

+arg: arg is input and should be instantiated.

-arg: arg is output and can be initially uninstantiated; if the query succeeds, the arg is instantiated with the "output" of the query.

?arg: arg can be either input or output

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online help

```
?- help(lists).  
No help available for lists  
Yes  
?- apropos(lists).  
merge/3           Merge two sorted lists  
append/3          Concatenate lists  
Section 11-1       "lists: List Manipulation"  
Section 15-2-1    "lists"  
Yes  
?- help(append/3).  
append(?List1, ?List2, ?List3)  
Succeeds when List3 unifies with the concatenation of List1 and  
List2. The predicate can be used with any instantiation pattern  
(even three variables).
```

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examples

```
?- append([a,b],[c],X).  
X = [a, b, c]  
  
Yes  
?- append(X,[c],[a,b,c]).  
X = [a, b]  
  
Yes  
?- append([a,b],[c],[a,b,d]).  
  
No
```

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more examples

```
?- append([a,b],X,Y).
X = _G187
Y = [a, b|_G187]
Yes
?- append(X,Y,Z).
X = []
Y = _G181
Z = _G181 ;

X = [_G262]
Y = _G181
Z = [_G262|_G181] ;

X = [_G262, _G268]
Y = _G181
Z = [_G262, _G268|_G181]
```

append is an example of a reversible or **steadfast** predicate (Richard O'Keefe)

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reversible programming

- good predicates are steadfast
- they give correct answers even if unusual values are supplied
 - e. g. variables for inputs, constants for outputs
- non-steadfast predicates require specific arguments to be instantiated (input) or variables (output)

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unification

- Prolog matches terms by *unifying* them, i.e. applying a most general unifier to them
- it instantiates variables as little as possible to make them match, e.g.

?- X = f(Y,b,Z), X = f(a,V,W).

X = f(a, b, _G182)

Y = a

Z = _G182

V = b

W = _G182

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family relations example

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family relations

- the database:

rules

```
parent(Parent, Child) :- mother(Parent, Child).
```

```
parent(Parent, Child) :- father(Parent, Child).
```

facts

```
father('George', 'Elizabeth'). father('George', 'Margaret').
```

```
mother('Mary', 'Elizabeth'). mother('Mary', 'Margaret').
```

- Note encoding of disjunction

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finding all solutions

```
| ?- parent(Parent, Child).
```

```
Parent = 'Mary',
```

```
Child = 'Elizabeth' ;
```

```
Parent = 'Mary',
```

```
Child = 'Margaret' ;
```

```
Parent = 'George',
```

```
Child = 'Elizabeth' ;
```

```
Parent = 'George',
```

```
Child = 'Margaret' ;
```

```
no
```

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how prolog finds solutions

```
trace] ?-
  parent(Parent, Child1),
  parent(Parent, Child2),
  not(Child1 = Child2).
Call: (8) parent(_G313,
_G314) ? creep
Exit: (9) mother(_G313,
_G314) ? creep
Exit: (9) mother('Mary',
'Elizabeth') ? creep
Exit: (8) parent('Mary',
'Elizabeth') ? creep
Call: (8) parent('Mary', _G317)
? creep
Call: (9) mother('Mary',
_G317) ? creep
Exit: (9) mother('Mary',
'Elizabeth') ? creep
Exit: (8) parent('Mary',
'Elizabeth') ? creep
Redo: (9) mother('Mary',
_G317) ? creep
Exit: (9) mother('Mary',
'Margaret') ? creep
Exit: (8) parent('Mary',
'Margaret') ? creep
Parent = 'Mary'
Child1 = 'Elizabeth'
Child2 = 'Margaret'
```

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Prolog's query answering process

- a query is a conjunction of terms
- answer to the query is *yes* if all terms *succeed*
- A term in a query *succeeds* if
 - it matches a fact in the database *or*
 - it matches the head of a rule whose body succeeds
- the substitution used to unify the term and the fact/head is applied to the rest of the query
- works on query terms in **left to right order**; databases facts/rules that match are tried in **top to bottom** order

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recursion examples

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generating permutations

- A permutation P of a list L is a list whose first is some element E of L and whose rest is a permutation of L with E removed.
- [] is a permutation of []
- In Prolog:

```
permutation([],[]).  
permutation(L,[E|PR]) :- select(E,L,R),  
                           permutation(R,PR).
```

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selecting an element from a list

- To select an element from a list, can either select the first leaving the rest, or select some element from the rest and leaving the first plus the unselected elements from the rest.
- In Prolog:
`select(X,[X|R],R).`
`select(X,[Y|R],[Y|RS]):- select(X,R,RS).`

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sorting by the definition

- Find a permutation that is ordered
`sort(L,P):- permutation(L,P),`
`ordered(P).`
`ordered([]).`
`ordered([E]).`
`ordered([E1,E2|R]) :- E1 <= E2,`
`ordered([E2|R]).`
- an example of “generate and test”

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reverse

- reverse(L,RL) holds if RL is a list with the components of L reversed
- ordinary recursive definition

```
reverse([],[]).
reverse([F|R],RL):- reverse(R,RR),
    append(RR, [F], RL).
append([],L,L).
append([F|R],L,[F|RL]):-
    append(R,L,RL).
```

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reverse

- Tail recursive definition:

```
reverse(L,RL):- reverse(L,[],RL).
reverse([],Acc,Acc).
reverse([F|R],Acc,RL):-
    reverse(R,[F|Acc],RL).
```
- recursive call is last thing done
- can avoid saving calls on stack

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solving a logic puzzle with Prolog

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the zebra puzzle

1. There are 5 houses, occupied by politically-incorrect gentlemen of 5 different nationalities, who all have different coloured houses, keep different pets, drink different drinks, and smoke different (now-extinct) brands of cigarettes.
2. The Englishman lives in a red house.
3. The Spaniard keeps a dog.
4. The owner of the green house drinks coffee.
- ...
6. The ivory house is just to the left of the green house.
- ...
11. The Chesterfields smoker lives next to a house with a fox.

Who owns the zebra and who drinks water?

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Prolog implementation

- represent the 5 houses by a structure of 5 terms
house(Colour, Nationality, Pet, Drink, Cigarettes)
- create a partial structure using variables, to be filled by the solution process
- specify constraints to instantiate variables

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house building

```
makehouses(0,[]).
```

```
makehouses(N,[house(Col, Nat, Pet, Drk, Cig)|List])  
    :- N>0, N1 is N - 1, makehouses(N1,List).
```

or more cleanly with anonymous variables:

```
makehouses(N,[house(_,_,_,_,_)|List])  
    :- N>0, N1 is N - 1, makehouses(N1,List).
```

Why is this equivalent? (See p. 159.)

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the empty houses

?- makehouses(5, List).

```
List = [house(_G233, _G234, _G235, _G236, _G237),
        house(_G245, _G246, _G247, _G248, _G249),
        house(_G257, _G258, _G259, _G260, _G261),
        house(_G269, _G270, _G271, _G272, _G273),
        house(_G281, _G282, _G283, _G284, _G285)]
```

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constraints

- The Englishman lives in a red house.
`house(red, englishman, _, _, _) on List,`
- The Spaniard keeps a dog.
`house(_, spaniard, dog, _, _) on List,`
- The owner of the green house drinks coffee.
`house(green, _, _, coffee, _) on List`
- The ivory house is just to the left of the green house
`sublist2([house(ivory, _, _, _, _),
 house(green, _, _, _, _)], List),`
- The Chesterfields smoker lives next to a house with a fox.
`nextto(house(_, _, _, _, chesterfields),
 house(_, _, fox, _, _), List),`

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defining the on operator

- `on` is a user-defined infix operator that is a version of `member/2`
- `:- op(100,xfy,on).`
`X on List :- member(X,List).`
amounts to
`X on [X|_].`
`X on [_|R] :- X on R.`

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predicates for defining constraints

- “just to the left of”? “lives next to”?
- define `sublist(S,L)`
`sublist2([S1, S2], [S1, S2 | _]) .`
`sublist2(S, [_ | T]) :- sublist2(S, T).`
- define `nextto` predicate
`nextto(H1, H2, L) :- sublist2([H1, H2], L).`
`nextto(H1, H2 ,L) :- sublist2([H2, H1], L).`

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translating the constraints

- The ivory house is just to the left of the green house
`sublist2([house(ivy, _, _, _),
 house(green, _, _, _)], List),`
- The Chesterfields smoker lives next to a house with a fox.
`nextto(house(_, _, _, chesterfields),
 house(_, _, fox, _), List),`

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looking for the zebra

- Who owns the zebra and who drinks water?
`find(ZebraOwner, WaterDrinker) :-
 makehouses(5, List),
 house(red, englishman, _, _, _) on List,
 ... % all other constraints
 house(_, WaterDrinker, _, water, _) on List,
 house(_, ZebraOwner, zebra, _, _) on List.`
- solution is generated and queried in the same clause
- neither water or zebra are mentioned in the constraints

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solving the puzzle

?- [zebra].

% zebra compiled 0.00 sec, 5,360 bytes

Yes

?- find(ZebraOwner, WaterDrinker).

ZebraOwner = japanese

WaterDrinker = norwegian ;

No

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how Prolog finds solution

After first 8 constraints:

List = [

house(red, englishman, snail, _G251, old_gold),

house(green, spaniard, dog, coffee, _G264),

house(ivory, ukrainian, _G274, tea, _G276),

house(green, _G285, _G286, _G287, _G288),

house(yellow, _G297, _G298, _G299, kools)]

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how Prolog solves the puzzle

Then need to satisfy “the owner of the third house drinks milk”, i.e.

```
List = [_ , _ , house( _ , _ , _ , milk, _ ), _ , _ ],
```

Can't be done with current instantiation of List. So Prolog will **backtrack** and find another.

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how Prolog solves the puzzle

The unique complete solution is

```
L = [  
house(yellow, norwegian, fox, water, kools),  
house(blue, ukrainian, horse, tea, chesterfields),  
house(red, englishman, snail, milk, old_gold),  
house(ivory, spaniard, dog, orange,  
lucky_strike),  
house(green, japanese, zebra, coffee,  
parliaments)]
```

See course web page for code of the example.

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