

Dynamic Scheduling

- In Simple pipelines, instructions are issued in order.
- If an instruction stalls, all instructions after it are stalled too (could be O.K. to execute).
- DIV.D F0,F2,F4
- ADD.D F10,F0,F8
- SUB.D F12,F8,F14 

Dynamic Scheduling

- Rearrange order of instructions to reduce stalls while maintaining data flow
- Advantages:
 - Compiler doesn't need to have knowledge of microarchitecture
 - Handles cases where dependencies are unknown at compile time
- Disadvantage:
 - Substantial increase in hardware complexity
 - Complicates exceptions

Dynamic Scheduling

- Rearrange order of instructions to reduce stalls while maintaining data flow
- Instructions are **issued** in program order
- But, the instruction begins execution as soon as its operands are ready
- Out of order execution → out of order completion
- DIV.D F0,F2,F4
- ADD.D F6,F0,F8 Antidependence
- SUB.D F8,F10,F14 Output Dependence
- MUL.D F6,F10,F8

Dynamic Scheduling

- To allow dynamic scheduling, split the ID stage in the simple MIPS pipeline into 2 stages
 - Issue: Decode and check for structural hazards
 - Read operand: wait for data hazard → read operand
- Instruction fetch stage before issue, and execution starts after read operand.
- Instructions pass through the issue stage in order, they can be delayed or pass each other at the read operand stage.

Dynamic Scheduling

- Major complication for exception handling.
- Must preserve the exception behavior as if the instructions are executed in the program order.
- May delay notification until the processor knows the instruction is the next one completed.
- Imprecise exception may occur
 - Later instructions (in program order) may have been completed already.
 - Earlier instructions may have not been completed

Register Renaming

- Example:

DIV.D	F0,F2,F4	
ADD.D	F6 ,F0, F8	antidependence
S.D	F6,0(R1)	
SUB.D	F8 ,F10,F14	antidependence
MUL.D	F6 ,F10,F8	

+ name dependence with F6

Register Renaming

- Example:

DIV.D F0,F2,F4	DIV.D F0,F2,F4
ADD.D F6 ,F0, F8	ADD.D S ,F0,F8
S.D F6,0(R1)	S.D S ,0(R1)
SUB.D F8 ,F10,F14	SUB.D T ,F10,F14
MUL.D F6 ,F10,F8	MUL.D F6 ,F10, T

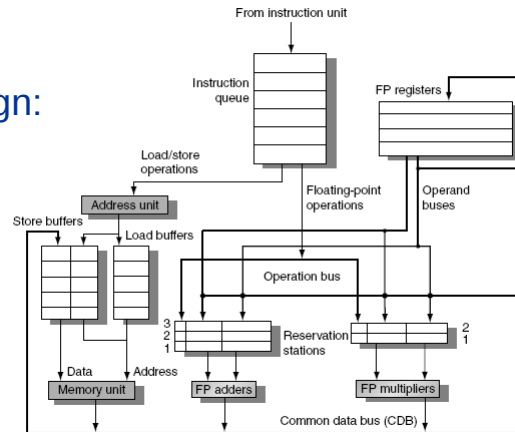
- Now only RAW hazards remain, which can be strictly ordered

Register Renaming

- Register renaming is provided by reservation stations (RS)
 - Contains:
 - The instruction
 - Buffered operand values (when available)
 - Reservation station number of instruction providing the operand values
 - RS fetches and buffers an operand as soon as it becomes available (not necessarily involving register file)
 - Pending instructions designate the RS to which they will send their output
 - Result values broadcast on a result bus, called the common data bus (CDB)
 - Only the last output updates the register file
 - As instructions are issued, the register specifiers are renamed with the reservation station
 - May be more reservation stations than registers

Tomasulo's Algorithm

- Load and store buffers
 - Contain data and addresses, act like reservation stations
- Top-level design:



Tomasulo's Algorithm

- Three Steps:
 - Issue
 - Get next instruction from FIFO queue
 - If available RS, issue the instruction to the RS with operand values if available
 - If no RS, stall the instruction issue
 - Execute
 - When operand becomes available, store it in any reservation stations waiting for it
 - When all operands are ready, start execution
 - Loads and store maintained in program order through effective address ****
 - No instruction allowed to initiate execution until all branches that proceed it in program order have completed
 - Write result
 - Write result on CDB into reservation stations and store buffers
 - (Stores must wait until address and value are received)

Tomasulo's Algorithm

- Op:** Operation to perform in the unit (e.g., + or −)
- Vj, Vk:** Value of Source operands
- Store buffers has V field, result to be stored
- Qj, Qk:** Reservation stations producing source registers (value to be written)
- Note: $Qj, Qk=0 \rightarrow$ ready
 - Store buffers only have Qj for RS producing result
- A:** Used to hold info for the load store (initially immediate, then effective address)
- Busy:** Indicates reservation station or FU is busy
- Register result status— **Qi** indicates which functional unit will write each register, 0 means no write to this register

Example

Instruction		Instruction status		
		Issue	Execute	Write Result
L.D	F6, 32(R2)	✓	✓	✓
L.D	F2, 44(R3)	✓	✓	
MUL.D	F0, F2, F4	✓		
SUB.D	F8, F2, F6	✓		
DIV.D	F10, F0, F6	✓		
ADD.D	F6, F8, F2	✓		

Reservation stations							
Name	Busy	Op	Vj	Vk	Qj	Qk	A
Load1	No						
Load2	Yes	Load					44 + Regs[R3]
Add1	Yes	SUB		Mem[32 + Regs[R2]]	Load2		
Add2	Yes	ADD			Add1	Load2	
Add3	No						
Mult1	Yes	MUL		Regs[F4]	Load2		
Mult2	Yes	DIV		Mem[32 + Regs[R2]]	Mult1		

Register status								
Field	F0	F2	F4	F6	F8	F10	F12	... F30
Qi	Mult1	Load2		Add2	Add1	Mult2		

Example

Instruction status							
Instruction				Issue	Execute	Write Result	
L.D F6,32(R2)				✓	✓	✓	
L.D F2,44(R3)				✓	✓		
MUL.D F0,F2,F4				✓			
SUB.D F8,F2,F6				✓			
DIV.D F10,F0,F6				✓			
ADD.D F6,F8,F2				✓			

Reservation stations							
Name	Busy	Op	Vj	Vk	Qj	Qk	A
Load1	No						
Load2	Yes	Load					44 + Regs[R3]
Add1	Yes	SUB		Mem[32 + Regs[R2]]	Load2		
Add2	Yes	ADD			Add1	Load2	
Add3	No						
Mult1	Yes	MUL		Regs[F4]	Load2		
Mult2	Yes	DIV		Mem[32 + Regs[R2]]	Mult1		

Register status							
Field	F0	F2	F4	F6	F8	F10	F12 ... F30
Qi	Mult1	Load2		Add2	Add1	Mult2	

Example

Instruction status							
Instruction				Issue	Execute	Write Result	
L.D F6,32(R2)				✓	✓	✓	
L.D F2,44(R3)				✓	✓		
MUL.D F0,F2,F4				✓			
SUB.D F8,F2,F6				✓			
DIV.D F10,F0,F6				✓			
ADD.D F6,F8,F2				✓			

Reservation stations							
Name	Busy	Op	Vj	Vk	Qj	Qk	A
Load1	No						
Load2	Yes	Load					44 + Regs[R3]
Add1	Yes	SUB		Mem[32 + Regs[R2]]	Load2		
Add2	Yes	ADD			Add1	Load2	
Add3	No						
Mult1	Yes	MUL		Regs[F4]	Load2		
Mult2	Yes	DIV		Mem[32 + Regs[R2]]	Mult1		

Register status							
Field	F0	F2	F4	F6	F8	F10	F12 ... F30
Qi	Mult1	Load2		Add2	Add1	Mult2	

Example

Instruction status							
Instruction	Issue			Execute		Write Result	
L.D F6,32(R2)	✓			✓		✓	
L.D F2,44(R3)	✓			✓			
MUL.D F0,F2,F4	✓						
SUB.D F8,F2,F6	✓						
DIV.D F10,F0,F6	✓						
ADD.D F6,F8,F2	✓						

Reservation stations							
Name	Busy	Op	Vj	Vk	Qj	Qk	A
Load1	No						
Load2	Yes	Load					44 + Regs[R3]
Add1	Yes	SUB		Mem[32 + Regs[R2]]	Load2		
Add2	Yes	ADD			Add1	Load2	
Add3	No						
Mult1	Yes	MUL		Regs[F4]	Load2		
Mult2	Yes	DIV		Mem[32 + Regs[R2]]	Mult1		

Register status							
Field	F0	F2	F4	F6	F8	F10	F12 ... F30
Qi	Mult1	Load2		Add2	Add1	Mult2	

Example

Instruction status							
Instruction	Issue			Execute		Write Result	
L.D F6,32(R2)	✓			✓		✓	
L.D F2,44(R3)	✓			✓			
MUL.D F0,F2,F4	✓						
SUB.D F8,F2,F6	✓						
DIV.D F10,F0,F6	✓						
ADD.D F6,F8,F2	✓						

Reservation stations							
Name	Busy	Op	Vj	Vk	Qj	Qk	A
Load1	No						
Load2	Yes	Load					44 + Regs[R3]
Add1	Yes	SUB		Mem[32 + Regs[R2]]	Load2		
Add2	Yes	ADD			Add1	Load2	
Add3	No						
Mult1	Yes	MUL		Regs[F4]	Load2		
Mult2	Yes	DIV		Mem[32 + Regs[R2]]	Mult1		

Register status							
Field	F0	F2	F4	F6	F8	F10	F12 ... F30
Qi	Mult1	Load2		Add2	Add1	Mult2	

Dealing with WAR

Instruction			Instruction status	
L.D	F6, 32(R2)			
L.D	F2, 44(R3)			
MUL.D	F0, F2, F4			
SUB.D	F8, F2, F6			
DIV.D	F10, F0, F6			
ADD.D	F6, F8, F2			

Name	Busy	Op
Load1	No	
Load2	Yes	Load
Add1	Yes	SUB
Add2	Yes	ADD
Add3	No	
Mult1	Yes	MUL
Mult2	Yes	DIV

Register status							
Field	F0	F2	F4	F6	F8	F10	F12 ... F30
Qi	Mult1	Load2		Add2	Add1	Mult2	

The processor issues both DIV and ADD although there is a WAR hazard.

If F6 is ready when DIV is issued, its value is read and stored in the RS (ADD may change it that is O.K.)

If not ready, RS will read it from the FU producing it, again ADD may change F6 since we will read it from the FU not F6

Instruction stream

Instruction status:

Instruction status:

Instruction	j	k	Issue	Comp	Result
LD	F6	34+	R2		
LD	F2	45+	R3		
MULTD	F0	F2	F4		
SUBD	F8	F6	F2		
DIVD	F10	F0	F6		
ADD	F6	F8	F2		

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

3 Load/Buffers

3 Load/Buffers

Reservation Stations:

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
	Mult2	No					

Annotations:

- 3 FP Adder R.S.
- 2 FP Mult R.S.
- FU count down

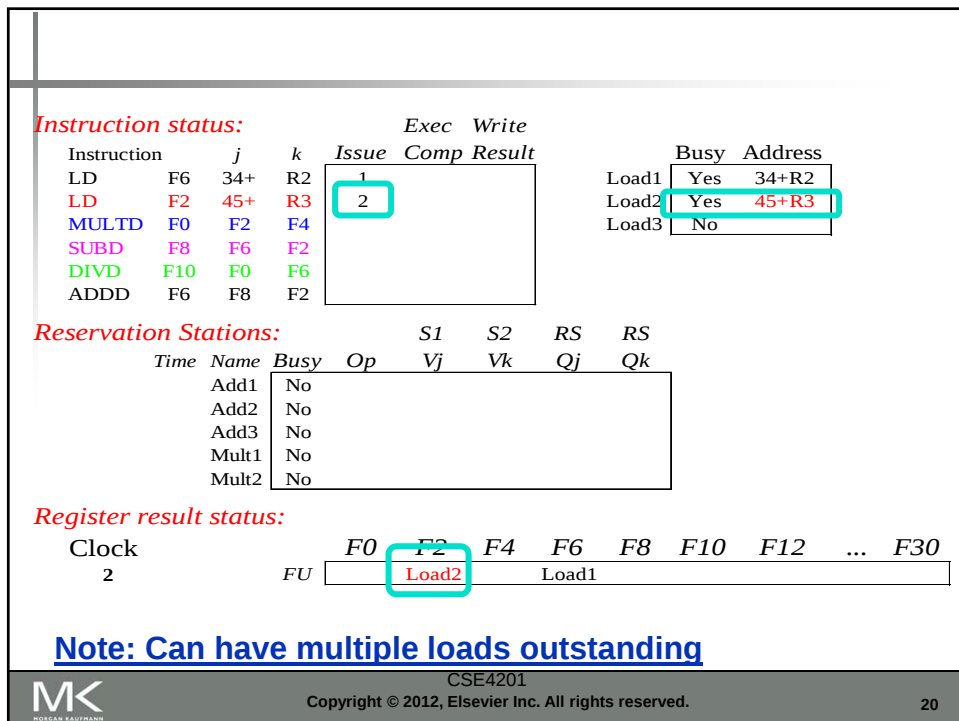
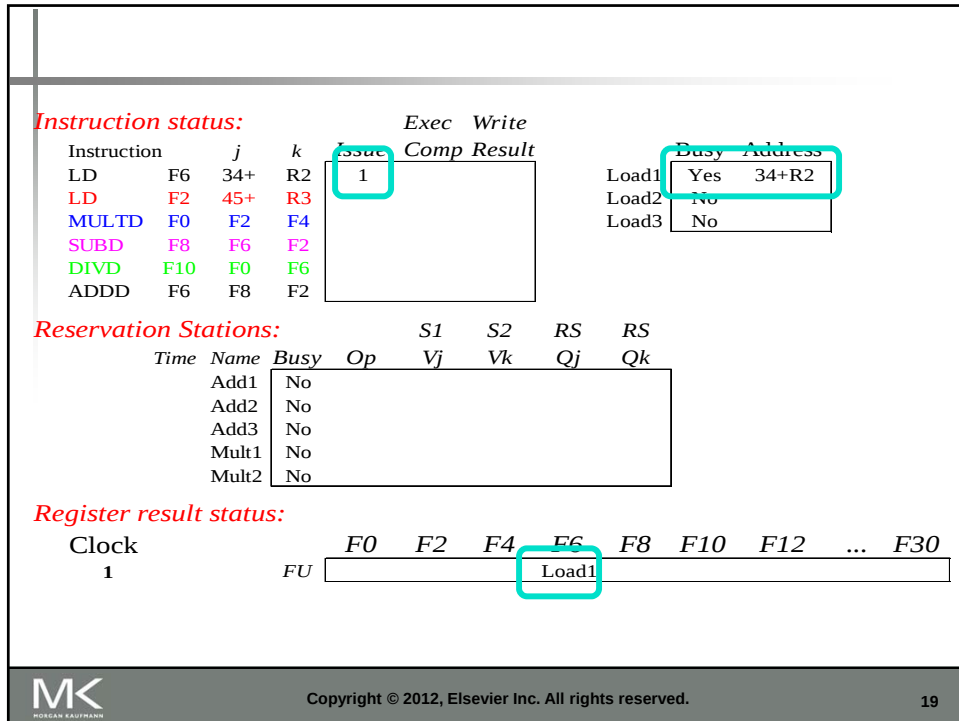
FU count down

3 FP Adder R.S.
2 FP Mult R.S.

Register result status:

Clock							
	F0	F2	F4	F6	F8	F10	F12 ... F30
FU							

Clock cycle counter



Instruction status:

Instruction	j	k	Issue	Exec		Write	Busy	Address
				Comp	Result			
LD	F6	34+	R2	1	3		Load1	Yes 34+R2
LD	F2	45+	R3	2			Load2	Yes 45+R3
MULTD	F0	F2	F4	3			Load3	No
SUBD	F8	F6	F2					
DIVD	F10	F0	F6					
ADDD	F6	F8	F2					

Reservation Stations:

Time	Name	Busy	Op	S1		S2		RS	
				Vj	Vk	Qj	Qk		
Add1		No							
Add2		No							
Add3		No							
Mult1		Yes	MULTD		R(F4)		Load2		
Mult2		No							

Register result status:

Clock		F0	F2	F4	F6	F8	F10	F12	...	F30
3	FU	Mult1	Load2		Load1					

- Note: registers names are removed ("renamed") in Reservation Stations; MULT issued
- Load1 completing; who is waiting for Load1?

Instruction status:

Instruction	j	k	Issue	Exec		Write	Busy	Address
				Comp	Result			
LD	F6	34+	R2	1	3	4	Load1	No
LD	F2	45+	R3	2	4		Load2	Yes 45+R3
MULTD	F0	F2	F4	3			Load3	No
SUBD	F8	F6	F2	4				
DIVD	F10	F0	F6					
ADDD	F6	F8	F2					

Reservation Stations:

Time	Name	Busy	Op	S1		S2		RS	
				Vi	Vk	Qj	Qk		
Add1		Yes	SUBD	M(A1)			Load2		
Add2		No							
Add3		No							
Mult1		Yes	MULTD		R(F4)		Load2		
Mult2		No							

Register result status:

Clock		F0	F2	F4	F6	F8	F10	F12	...	F30
4	FU	Mult1	Load2		M(A1)	Add1				

- Load1 completing; what is waiting for Load1?

Instruction status:

Instruction	j	k	Issue	Exec Write		Busy	Address
				Comp	Result		
LD	F6	34+	R2	1	3	4	Load1
LD	F2	45+	R3	2	4	5	Load2
MULTD	F0	F2	F4	3			Load3
SUBD	F8	F6	F2	4			
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2				

Reservation Stations:

Time	Name	Busy	Op	Vj	Vk	Qj	Qk
2	Add1	Yes	SUBD	M(A1)	M(A2)		
	Add2	No					
	Add3	No					
10	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD	M(A1)	Mult1		

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
5	FU	Mult1	M(A2)		M(A1)	Add1	Mult2		

- Timer starts down for Add1, Mult1

Instruction status:

Instruction	j	k	Issue	Exec Write		Busy	Address
				Comp	Result		
LD	F6	34+	R2	1	3	4	Load1
LD	F2	45+	R3	2	4	5	Load2
MULTD	F0	F2	F4	3			Load3
SUBD	F8	F6	F2	4			
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6			

Reservation Stations:

Time	Name	Busy	Op	Vj	Vk	Qj	Qk
1	Add1	Yes	SUBD	M(A1)	M(A2)		
	Add2	Yes	ADDD		M(A2)	Add1	
	Add3	No					
9	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD	M(A1)	Mult1		

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
6	FU	Mult1	M(A2)		Add2	Add1	Mult2		

- Issue ADDD here despite name dependency on F6?

Instruction status:

Instruction	j	k	Issue	Exec Write		Busy	Address
				Comp	Result		
LD	F6	34+	R2	1	3	4	Load1
LD	F2	45+	R3	2	4	5	Load2
MULTD	F0	F2	F4	3			Load3
SUBD	F8	F6	F2	4	7		
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6			

Reservation Stations:

Time	Name	Busy	Op	S1 S2		RS RS	
				Vj	Vk	Qj	Qk
0	Add1	Yes	SUBD	M(A1)	M(A2)		
	Add2	Yes	ADDD		M(A2)	Add1	← waiting
	Add3	No					
8	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
7	FU	Mult1	M(A2)		Add2	Add1	Mult2		

- Add1 (SUBD) completing; what is waiting for it?

Instruction status:

Instruction	j	k	Issue	Exec Write		Busy	Address
				Comp	Result		
LD	F6	34+	R2	1	3	4	Load1
LD	F2	45+	R3	2	4	5	Load2
MULTD	F0	F2	F4	3			Load3
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6			

Reservation Stations:

Time	Name	Busy	Op	S1 S2		RS RS	
				Vj	Vk	Qj	Qk
	Add1	No					
2	Add2	Yes	ADDD	(M-M)	M(A2)		
	Add3	No					
7	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
8	FU	Mult1	M(A2)		Add2	(M-M)	Mult2		

Instruction status:

Instruction	<i>j</i>	<i>k</i>	<i>R2</i>	Exec		Write		Busy	Address
				Issue	Comp	Result			
LD	F6	34+	R2	1	3	4	Load1	No	
LD	F2	45+	R3	2	4	5	Load2	No	
MULTD	F0	F2	F4	3			Load3	No	
SUBD	F8	F6	F2	4	7	8			
DIVD	F10	F0	F6	5					
ADDD	F6	F8	F2	6					

Reservation Stations:

Time	Name	Busy	Op	S1		S2		RS	
				<i>Vj</i>	<i>Vk</i>	<i>Qj</i>	<i>Qk</i>		
	Add1	No							
1	Add2	Yes	ADDD	(M-M)	M(A2)				
	Add3	No							
6	Mult1	Yes	MULTD	M(A2)	R(F4)				
	Mult2	Yes	DIVD		M(A1)	Mult1			

Register result status:

Clock	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
9	FU								
	Mult1	M(A2)		Add2	(M-M)	Mult2			

Instruction status:

Instruction	<i>j</i>	<i>k</i>	<i>R2</i>	Exec		Write		Busy	Address
				Issue	Comp	Result			
LD	F6	34+	R2	1	3	4	Load1	No	
LD	F2	45+	R3	2	4	5	Load2	No	
MULTD	F0	F2	F4	3			Load3	No	
SUBD	F8	F6	F2	4	7	8			
DIVD	F10	F0	F6	5					
ADDD	F6	F8	F2	6	10				

Reservation Stations:

Time	Name	Busy	Op	S1		S2		RS	
				<i>Vj</i>	<i>Vk</i>	<i>Qj</i>	<i>Qk</i>		
	Add1	No							
0	Add2	Yes	ADDD	(M-M)	M(A2)				
	Add3	No							
5	Mult1	Yes	MULTD	M(A2)	R(F4)				
	Mult2	Yes	DIVD		M(A1)	Mult1			

Register result status:

Clock	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
10	FU								
	Mult1	M(A2)		Add2	(M-M)	Mult2			

- Add2 (ADDD) completing; what is waiting for it?

Instruction status:

Instruction	j	k	Exec Write			Busy	Address
			Issue	Comp	Result		
LD	F6	34+	R2	1	3	4	Load1
LD	F2	45+	R3	2	4	5	Load2
MULTD	F0	F2	F4	3			Load3
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6	10	11	

Reservation Stations:

Time	Name	Busy	Op	S1		S2		RS	
				Vj	Vk	Qj	Qk	RS	
	Add1	No							
	Add2	No							
	Add3	No							
4	Mult1	Yes	MULTD	M(A2)		R(F4)			
	Mult2	Yes	DIVD		M(A1)	Mult1			

Register result status:

Clock		F0 F2 F4 F6 F8 F10 F12 ... F30											
		FU Mult1 M(A2) (M-M+N) (M-M) Mult2											
11													

- Write result of ADDD here?
- All quick instructions complete in this cycle!

Instruction status:

Instruction	j	k	Exec Write			Busy	Address
			Issue	Comp	Result		
LD	F6	34+	R2	1	3	4	Load1
LD	F2	45+	R3	2	4	5	Load2
MULTD	F0	F2	F4	3			Load3
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6	10	11	

Reservation Stations:

Time	Name	Busy	Op	S1		S2		RS	
				Vj	Vk	Qj	Qk	RS	
	Add1	No							
	Add2	No							
	Add3	No							
3	Mult1	Yes	MULTD	M(A2)		R(F4)			
	Mult2	Yes	DIVD		M(A1)	Mult1			

Register result status:

Clock		F0 F2 F4 F6 F8 F10 F12 ... F30											
		FU Mult1 M(A2) (M-M+N) (M-M) Mult2											
12													

Instruction status:

Instruction	<i>j</i>	<i>k</i>	<i>Issue</i>	<i>Exec Write</i>		Busy	Address
				<i>Comp</i>	<i>Result</i>		
LD	F6	34+	R2	1	3	4	Load1 Load2 Load3
LD	F2	45+	R3	2	4	5	
MUL TD	F0	F2	F4	3			
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6	10	11	

Reservation Stations:

Time	Name	Busy	Op	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>
				<i>Vj</i>	<i>Vk</i>	<i>Qj</i>	<i>Qk</i>
	Add1	No					
	Add2	No					
	Add3	No					
2	Mult1	Yes	MUL TD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
13	FU	Mult1	M(A2)	(M-M+N	(M-M)	Mult2			

Instruction status:

Instruction	<i>j</i>	<i>k</i>	<i>Issue</i>	<i>Exec Write</i>		Busy	Address
				<i>Comp</i>	<i>Result</i>		
LD	F6	34+	R2	1	3	4	Load1 Load2 Load3
LD	F2	45+	R3	2	4	5	
MUL TD	F0	F2	F4	3			
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6	10	11	

Reservation Stations:

Time	Name	Busy	Op	<i>S1</i>	<i>S2</i>	<i>RS</i>	<i>RS</i>
				<i>Vj</i>	<i>Vk</i>	<i>Qj</i>	<i>Qk</i>
	Add1	No					
	Add2	No					
	Add3	No					
1	Mult1	Yes	MUL TD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	<i>F0</i>	<i>F2</i>	<i>F4</i>	<i>F6</i>	<i>F8</i>	<i>F10</i>	<i>F12</i>	...	<i>F30</i>
14	FU	Mult1	M(A2)	(M-M+N	(M-M)	Mult2			

Instruction status:

Instruction	j	k	Exec Write			Busy	Address
			Issue	Comp	Result		
LD	F6	34+	R2	1	3	4	Load1 Load2 Load3
LD	F2	45+	R3	2	4	5	
MULTD	F0	F2	F4	3	15		
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6	10	11	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
0	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
15	FU	Mult1	M(A2)		(M-M+N	(M-M)	Mult2		

- Mult1 (MULTD) completing; who is waiting for it?

Instruction status:

Instruction	j	k	Exec Write			Busy	Address
			Issue	Comp	Result		
LD	F6	34+	R2	1	3	4	Load1 Load2 Load3
LD	F2	45+	R3	2	4	5	
MULTD	F0	F2	F4	3	15	16	
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5			
ADDD	F6	F8	F2	6	10	11	

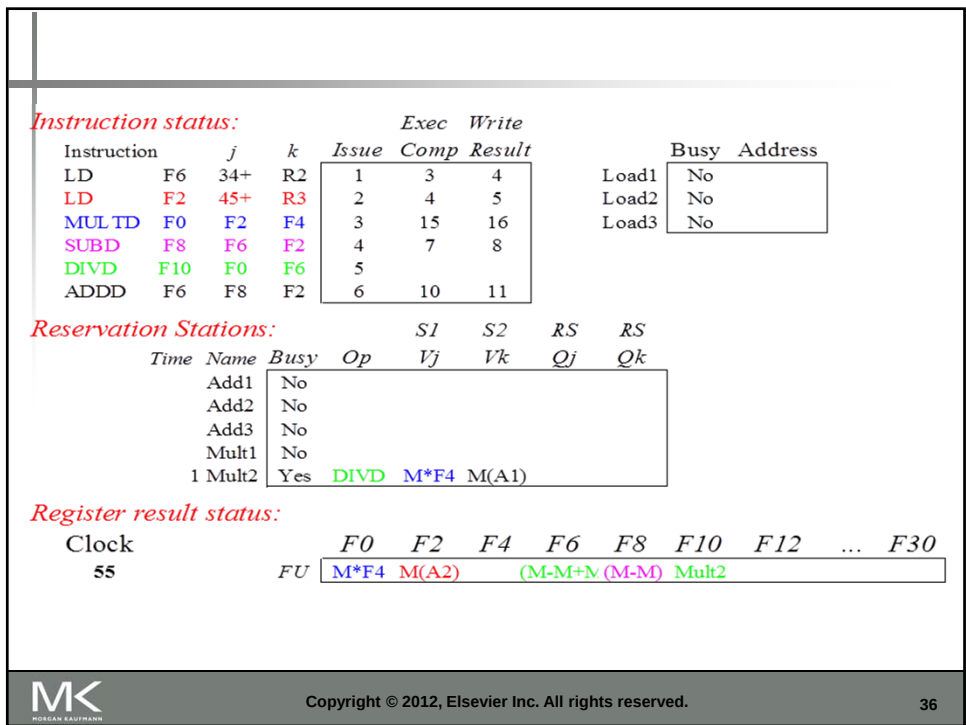
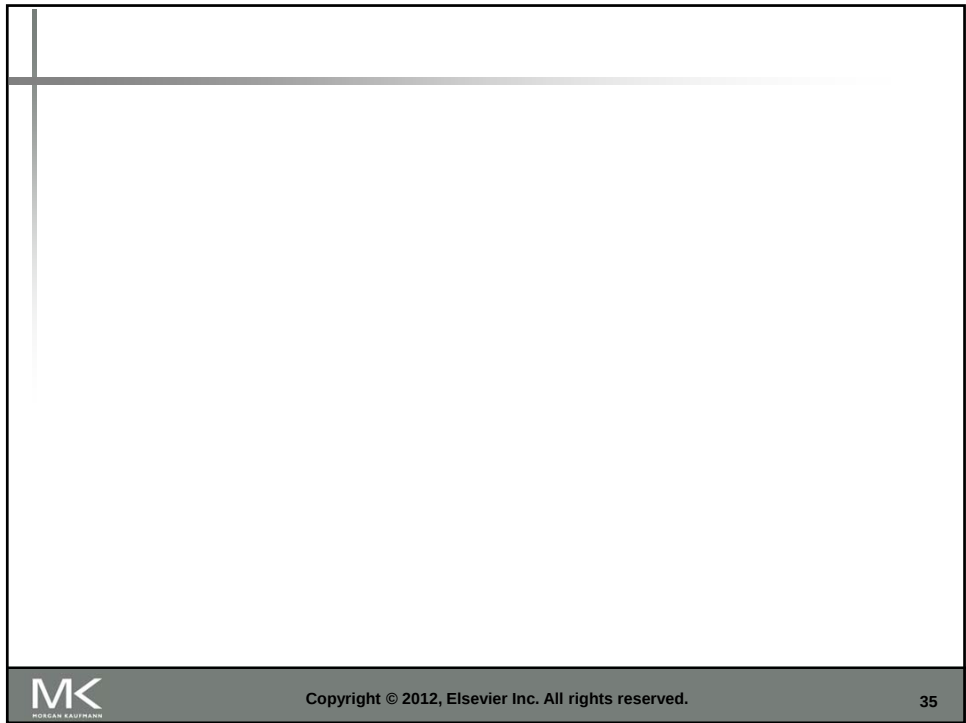
Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
40	Mult2	Yes	DIVD		M(A1)		

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
16	FU	M*A4	M(A2)		(M-M+N	(M-M)	Mult2		

- Just waiting for Mult2 (DIVD) to complete



Instruction status:

Instruction	j	k	Exec Write			Busy	Address
			Issue	Comp	Result		
LD	F6	34+	R2	1	3	4	Load1 Load2 Load3
LD	F2	45+	R3	2	4	5	
MULTD	F0	F2	F4	3	15	16	
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5	56		
ADDD	F6	F8	F2	6	10	11	

Reservation Stations:

Time	Name	Busy	Op	S1		S2		RS	
				Vj	Vk	Qj	Qk		
	Add1	No							
	Add2	No							
	Add3	No							
	Mult1	No							
0	Mult2	Yes	DIVD	M*F4	M(A1)				

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
56	M*F4	M(A2)		(M-M+N	(M-M)	Mult2			

Instruction status:

Instruction	j	k	Exec Write			Busy	Address
			Issue	Comp	Result		
LD	F6	34+	R2	1	3	4	Load1 Load2 Load3
LD	F2	45+	R3	2	4	5	
MULTD	F0	F2	F4	3	15	16	
SUBD	F8	F6	F2	4	7	8	
DIVD	F10	F0	F6	5	56	57	
ADDD	F6	F8	F2	6	10	11	

Reservation Stations:

Time	Name	Busy	Op	S1		S2		RS	
				Vj	Vk	Qj	Qk		
	Add1	No							
	Add2	No							
	Add3	No							
	Mult1	No							
	Mult2	Yes	DIVD	M*F4	M(A1)				

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
56	M*F4	M(A2)		(M-M+N	(M-M)	Result			

- Once again: In-order issue, out-of-order execution and out-of-order completion.

Tomasulo's Algorithm

- Load and stores could be done out of order provided they access different memory locations
- If they access same location, must preserve order (WAR, RAW, or WAW).
- If address calculation is done in program order, load/store can check if any uncompleted load/store share the same address
- Either wait or forward if possible.