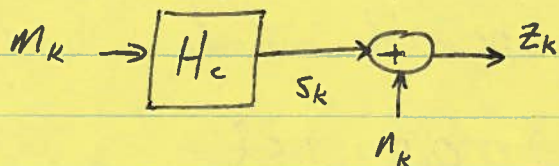


# L13 Sequence Detection

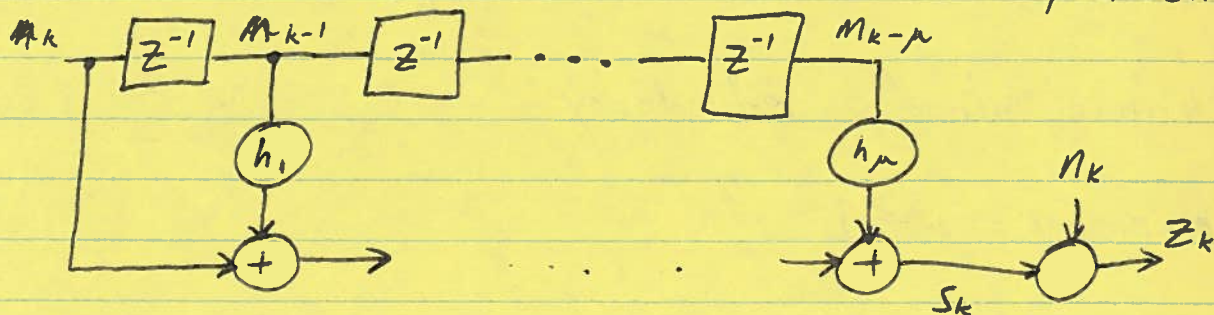
(1)

## 13.1 Channel Response

- after channel + electronics



- generally a system with memory ...  $\mu$  (# of previous symbols repeated remembered)



- the output is a convolution in discrete time

$$s_k = \sum_{l=0}^{\mu} h_l m_{k-l} = \sum_{l=0}^{\mu} h_{k-l} m_l$$

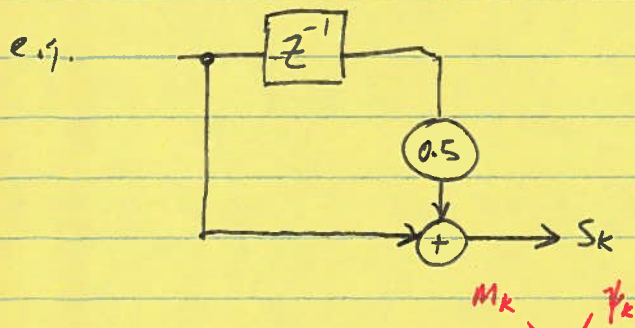
- effectively the output depends on...

$$s_k = f(m_k, \bar{\psi}_k)$$

current input  $\bar{\psi}_k = [m_{k-1}, m_{k-2}, \dots, m_{k-\mu}]$  current state of channel

- or entirely in terms of current and upcoming states

$$s_k = g(\bar{\psi}_k, \bar{\psi}_{k+1}) \leftarrow \text{o/p is fn. of a state transition}$$



- for alphabet  $\{0, 1\}$  of size  $M=2$

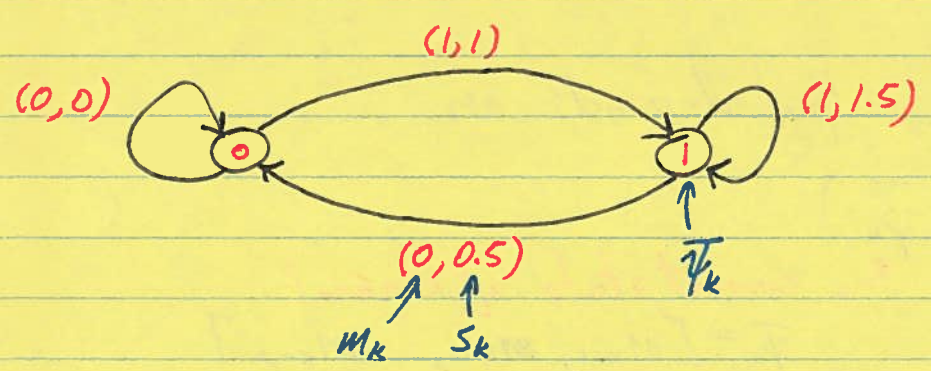
- possible outputs

$$S_k = \{S_k^{00}, S_k^{01}, S_k^{10}, S_k^{11}\} = \{0, 0.5, 1, 1.5\}$$

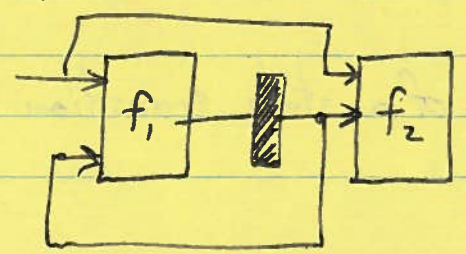
- 1-bit of info comes in + one of 4 levels comes out
- channel introduces redundancy
- memory =  $\mu = 1$

### 13.2 FSMs

- A generic representation of our channel is in the form of a **state transition diagram**

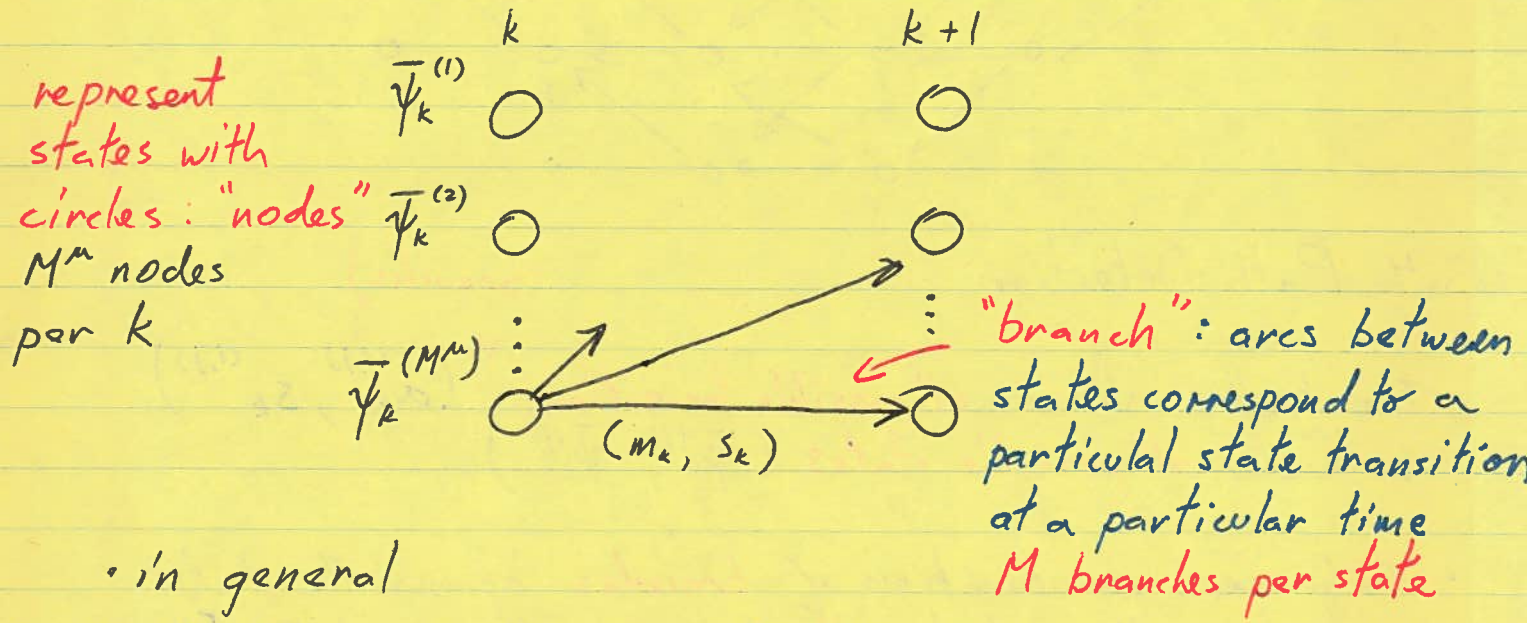


- which itself is a description of a FSM

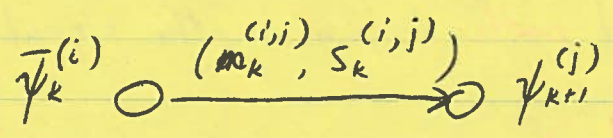


### 13.3 Trellis Diagram

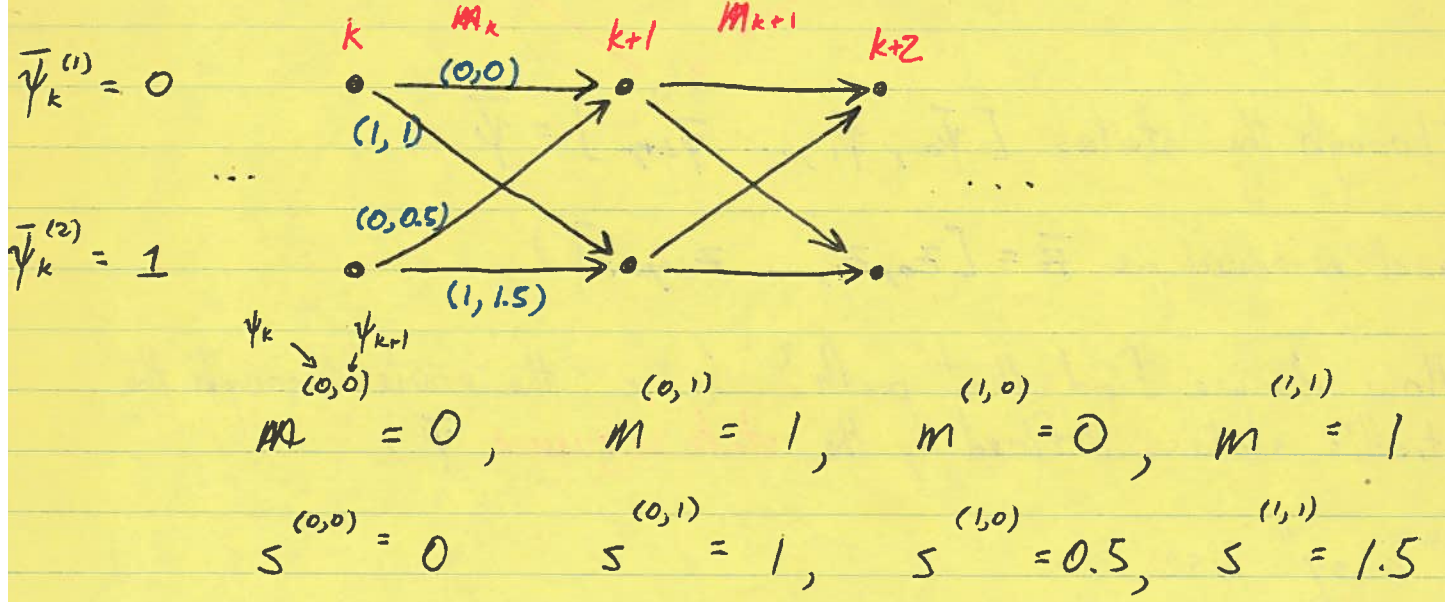
- Alternative to traditional state transition diagram
- Shows all possible progressions of states over time



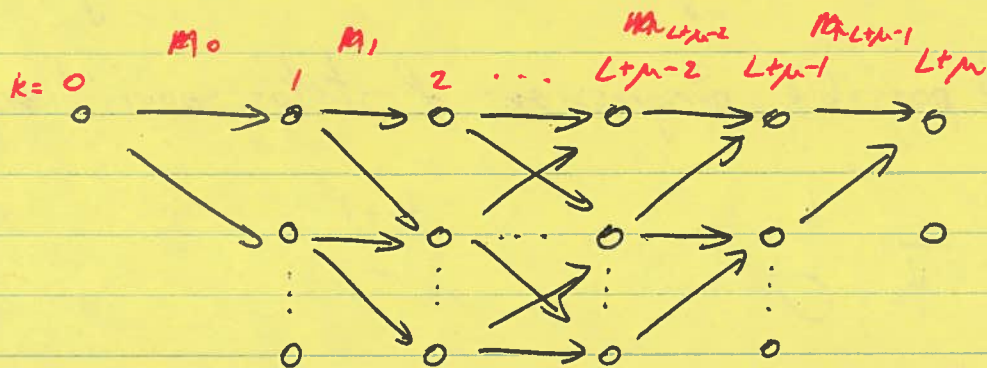
• in general



e.g. our simple channel in trellis form



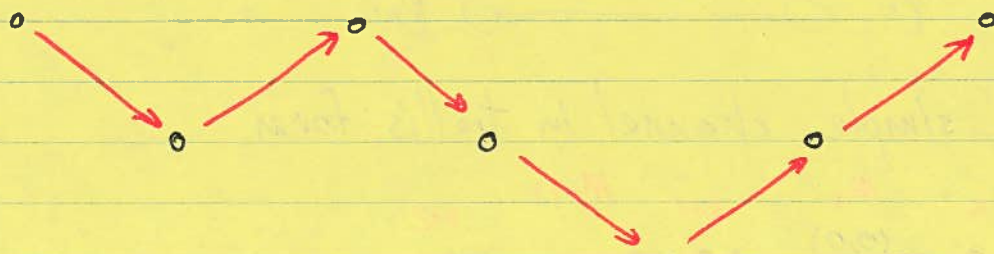
- bounds are often imposed by assuming that the sequence starts at some known state and ends at some known state



### 13.4 Path Detection

(branches)

- The trellis shows all possible transitions  $(a_k^{(i,j)}, s_k^{(i,j)})$  between all possible states  $(\bar{\psi}_k^{(i)}, \bar{\psi}_{k+1}^{(j)})$
- Only one combination of branches across the trellis constitutes the actual path taken by the signal  $\bar{m} = [m_0, m_1, \dots, m_{L+\mu}]$



through the states  $[\bar{\psi}_0, \bar{\psi}_1, \dots, \bar{\psi}_{L+\mu}] = \bar{\psi}$

(and received as  $\bar{z} = [z_0, z_1, \dots, z_{L+\mu-1}]$ )

- How do we find that path? (i.e. the route through the trellis states defined by the state sequence  $\bar{\psi}$ )
- "easy" use...

... MAP

$$\hat{\bar{\psi}} = \arg \max_{\bar{\psi}} p(\bar{\psi} | \bar{z})$$

$$p(\bar{\psi} | \bar{z}) = p(\bar{z} | \bar{\psi}) p(\bar{\psi})$$

$$p(\bar{z} | \bar{\psi}) = \prod_{k=0}^{L+\mu-1} p_{z_k}(z_k | \bar{\psi})$$

probability of sequence is product of probabilities of constituent transitions further,  $z_k$  depends only on adjacent (in time) states so...

$$p(\bar{\psi}) = p(\psi_0) \prod_{k=0}^{L+\mu-1} p(\bar{\psi}_{k+1} | \bar{\psi}_k)$$

product of probabilities of corresponding state transitions (and initial state)

$$p(\bar{z} | \bar{\psi}) = \prod_{k=0}^{L+\mu-1} p_z(z_k | \bar{\psi}_k, \bar{\psi}_{k+1})$$

$$\hat{\bar{\psi}} = \arg \max_{\bar{\psi}} \prod_{k=0}^{L+\mu-1} [p_z(z_k | \bar{\psi}_k, \bar{\psi}_{k+1}) p(\bar{\psi}_{k+1} | \bar{\psi}_k)]$$

for  $\bar{\psi}_k^{(i)} = i$  and  $\bar{\psi}_{k+1}^{(j)} = j$ , recall  $s_k = g(\bar{\psi}_k, \bar{\psi}_{k+1})$ ,

$$= \arg \max_{\bar{\psi}} \left\{ \prod_{k=0}^{L+\mu-1} p_z(z_k | s^{(i,j)}) \cdot p_{m_k}(m^{(i,j)}) \right\}$$

branch metrics

$p(\bar{\psi}_{k+1} | \bar{\psi}_k)$  depends on input  $m_k$

• the path value consists of  $L + \mu$  branch metrics

$$B_k(i, j) = p_z(z_k | s^{(i,j)}) \cdot p_{m_k}(m^{(i,j)})$$

↑  
branch metric & time k from states i to j

• for Gaussian noise

$$B_k(i, j) = p_N(z_k - s^{(i,j)}) \cdot p_{m_k}(m^{(i,j)})$$

$$= \frac{1}{\sqrt{2\pi\sigma_0^2}} e^{-\frac{|z_k - s^{(i,j)}|^2}{2\sigma_0^2}} \cdot p_{m_k}(m^{(i,j)})$$

$$-\log(B_k(i, j)) = \frac{|z_k - s^{(i,j)}|^2}{2\sigma_0^2} + \frac{1}{2} \log(2\pi\sigma_0^2) - \log[p_{m_k}(m^{(i,j)})]$$

$$B_k(i, j) \Big|_{\text{MAP}} = |z_k - s^{(i,j)}|^2 + \sigma_0^2 \log(2\pi\sigma_0^2) - 2\sigma_0^2 \log[p_{m_k}(m^{(i,j)})]$$

• for ML all  $m^{(i,j)}$  are equally likely hence  $p_{m_k}(m^{(i,j)})$  same for all branches

•  $\sigma_0^2$  same for all branches as well and hence these two components are just offsets which the maximum likelihood sequence detector (MLSD) can ignore

$$B_k(i, j) \Big|_{\text{ML}} = |z_k - s^{(i,j)}|^2$$

• because of our manipulations our path metric becomes

$$P \Big|_{\text{ML}} = \sum_{k=0}^{L+k-1} |z_k - s^{(i,j)}|^2 \quad : \text{sum of squares}$$

and... ↑ our "cost function"

$\hat{\psi} = \arg \min_{\bar{\psi}} P|_{ML} \leftarrow$  that is our ML path selection criterion is a **least squares** criterion

### 13.5 Viterbi Algorithm

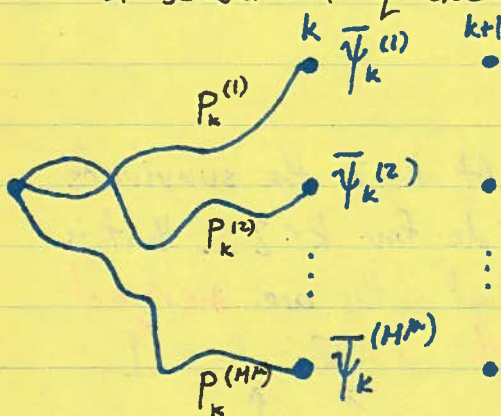
Finding the sequence that minimizes...

$P|_{ML} = \sum_{k=0}^{L+M-1} |z_k - s^{(i,j)}|^2$  ... is a tall order

- we have  $M^M$  states
- $M$  branches leaving & entering each state
- $L$  time intervals

therefore about  $(M^{M+1})^L$  possible sequences

- luckily this is amenable to minimization via **dynamic programming**  $\uparrow$  i.e. finding the min path metric
- a means of iteratively computing a sequence of steps from a selection of choices



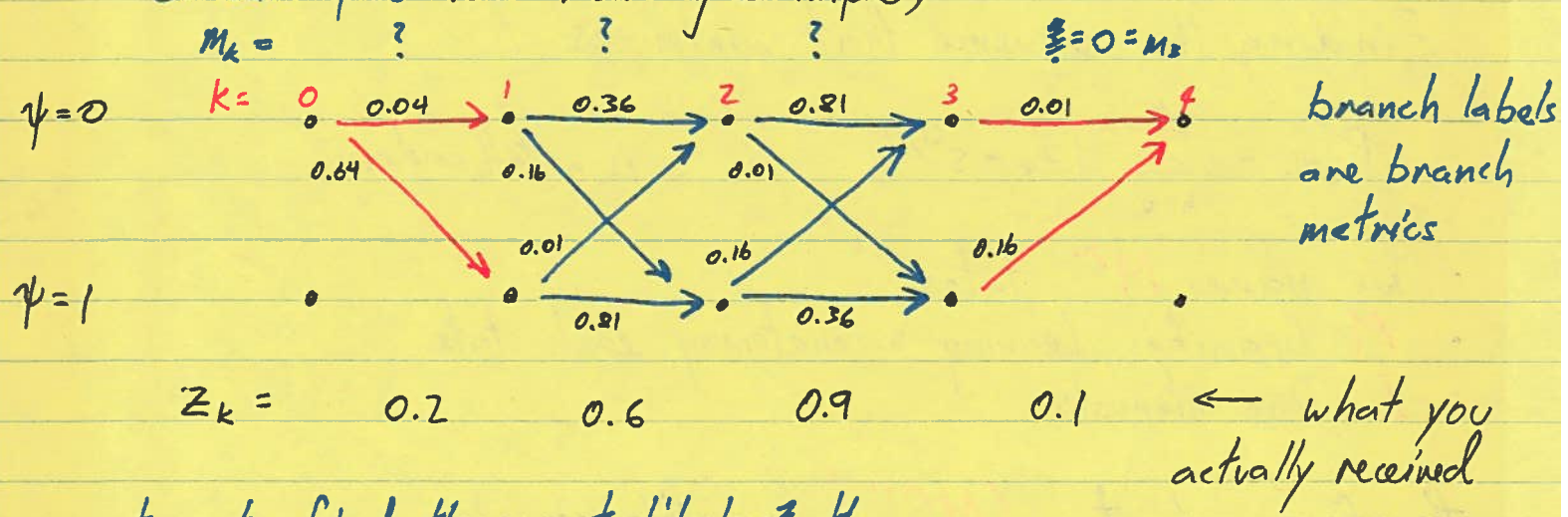
- ① identify optimal paths to states at time  $k, P_k$  (partial paths) with their **partial path metrics**
- ② calculate all **branch metrics** from  $k$  to  $k+1 : B_k(i,j)$  [i.e.  $M$  partial paths at each state  $\bar{\psi}_{k+1}^{(j)}$ ].  

$$P_{k+1}^{(j)} = \{ P_k^{(i)} + B_k(i,j) \}$$
 set of  $M$  partial paths into ea. state  $\bar{\psi}_{k+1}^{(j)}$
- ③ choose **survivor path** at each stage  

$$P_{k+1}^{(j)} = \min \{ P_k^{(i)} + B_k(i,j) \}$$

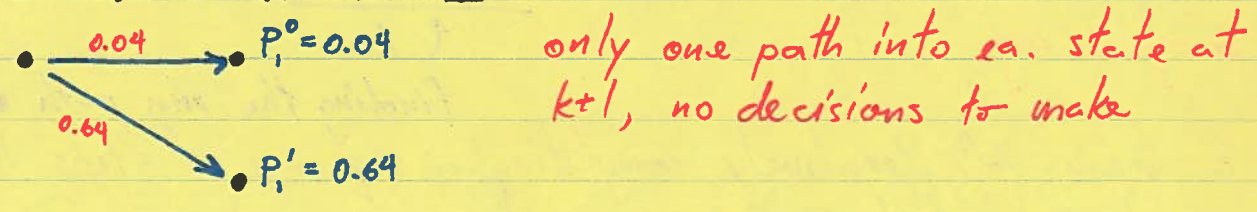
- in the case where the trellis starts and ends in a known state (i.e.  $m_{-1}$  is known as are  $m_L$  to  $m_{L+\mu}$ ) the remaining path (i.e. the **survivor sequence** at time  $L+\mu$ ) is the optimal path

- for example (our running example)

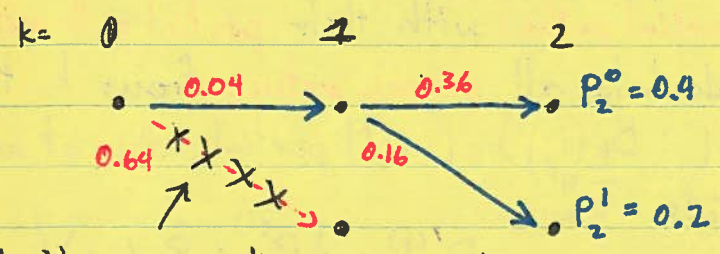


- try to find the most likely path...

1) survivors at  $k=1$



2) survivors at  $k=2$



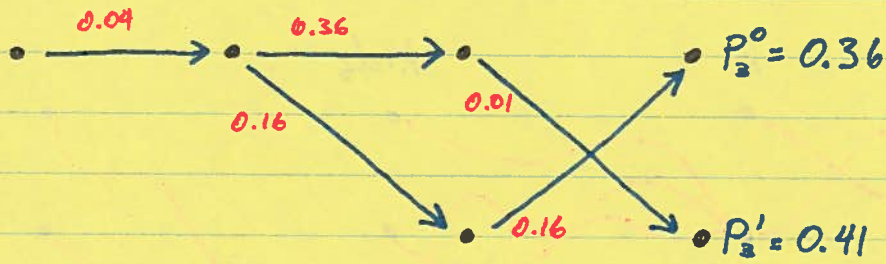
NOTE: At  $k=2$  the survivors all coincide for  $k < 2$  ... that is the partial paths are **merged** at **depth**  $d = 2 - 1 = 1$

$\nearrow$        $\uparrow$   
 $k=2$      $k=1$

don't even bother expanding this path from  $k=1$  as the  $B_1(i,j)$ 's from the other state are  $< 0.64$

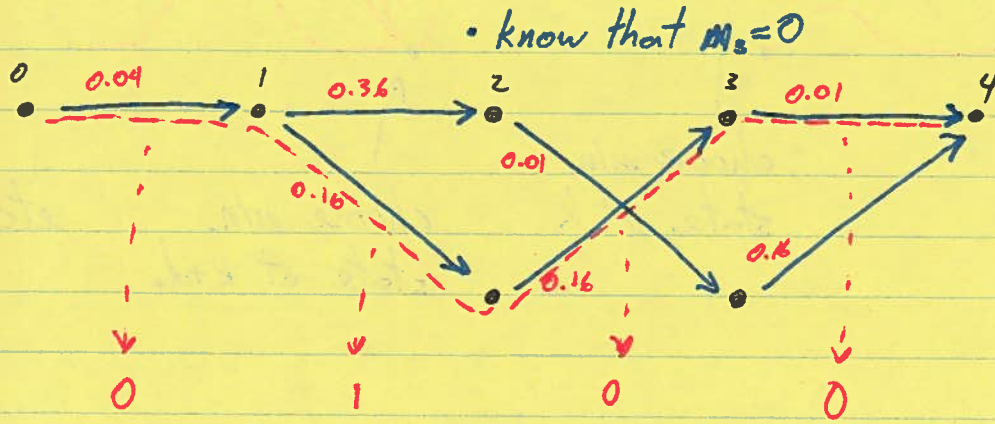


3) survivors at  $k=3$



$d=2$   
(merge depth)

4) survivors at  $k=4$



$d=3$

### parting thoughts

- having to wait for a whole sequence to finish can be quite restricting
- **path merges** are nice because they allow us to start with a new sequence, but they are not guaranteed to happen so we can't expect to wait around for them
- in practice you can always choose to employ some **truncation depth  $d_t$**

• at  $k$  find cheapest path and use that as your origin... repeat every  $\Delta t$  time steps

