

# Probabilistic Parsing

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# Outline

- 1 Context
- 2 PCFG as a Probabilistic Model
- 3 Computational Tasks
- 4 Probabilistic Inference
- 5 CNF and CYK
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# Probabilistic Context-Free Grammars

- N-gram and Hidden Markov models are linear
- Natural language syntax is not linear in structure
- Bayesian Networks is one way to capture structure
- PCFG, a derivative of CFG, is another way

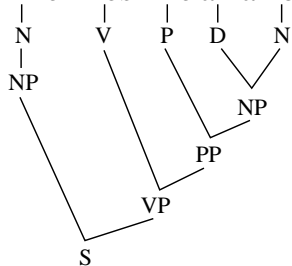
# Probabilistic Context-Free Grammars

- Existing CFG parsers are very good  
Example: CYK,  $O(n^3)$
- Can't apply to NL due to ambiguity

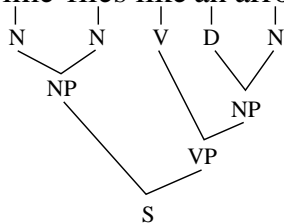
# Probabilistic Context-Free Grammars

Consider a sentence with two parse trees:

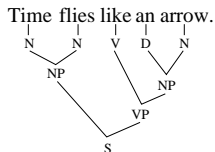
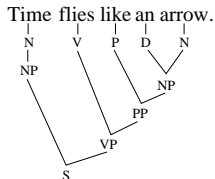
Time flies like an arrow.



Time flies like an arrow.



# Probabilistic Context-Free Grammars



These trees induce a CFG:

S	→	NP VP	VP	→	V NP	N	→	time	V	→	like
NP	→	N	VP	→	V PP	N	→	arrow	V	→	flies
NP	→	N N	PP	→	P NP	N	→	flies	P	→	like
NP	→	D N				D	→	an			

# Probabilistic Context-Free Grammars

- If we parse the same sentence using this CFG, we will obtain at least two different trees
- Need a way to assign a score to each tree

# PCFG as a Probabilistic Model

- Let's model derivations as stochastic processes

Consider the left-most derivation of the first tree:

$$\begin{aligned} \mathbf{S} &\rightarrow \mathbf{NP VP} \rightarrow \mathbf{N VP} \rightarrow \mathit{time VP} \rightarrow \mathit{time V PP} \\ &\rightarrow \mathit{time flies PP} \rightarrow \mathit{time flies P NP} \rightarrow \dots \end{aligned}$$

Each “ $\rightarrow$ ” involves a choice.



# PCFG as a Probabilistic Model

- Assign a probability  $P$  to each rule  $X \rightarrow \alpha$   
so that  $\sum_{i=1}^n P(X \rightarrow \alpha_i) = 1$  for each unique non-terminal  $X$

**S**  $\xrightarrow{1.0}$  **NP VP**

**NP**  $\xrightarrow{0.4}$  **N**

**NP**  $\xrightarrow{0.2}$  **NN**

**NP**  $\xrightarrow{0.4}$  **DN**

**V**  $\xrightarrow{0.3}$  *like*

**V**  $\xrightarrow{0.7}$  *flies*

**P**  $\xrightarrow{1.0}$  *like*

**VP**  $\xrightarrow{0.5}$  ...

This constitutes a PCFG.

# PCFG as a Probabilistic Model

- A **PCFG** is a CFG in which every production rule is associated with a probability.
- A PCFG is **proper** if its probability distribution is proper<sup>1</sup> over every subset of rules that have the same left-hand-side.
- A PCFG is **consistent** if its probability distribution is proper over the set of trees<sup>2</sup> it generates

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<sup>1</sup>i.e., adds up to 1, see previous slide

<sup>2</sup>or sentences, doesn't matter

# PCFG as a Probabilistic Model

- With rule probabilities, can calculate probability of a tree:

$$\begin{aligned}P(\text{TREE1}) &= P(\mathbf{N} \rightarrow \textit{time}) \times P(\mathbf{V} \rightarrow \textit{flies}) \times P(\mathbf{P} \rightarrow \textit{like}) \\ &\quad \times \dots \times P(\mathbf{S} \rightarrow \mathbf{NP VP}) \\ &= 0.0084\end{aligned}$$

$$P(\text{TREE2}) = 0.00036$$

- Now we have a winner

# PCFG as a Probabilistic Model

But where do the probabilities come from?

- Recall the Big Assignment
- Given set of parse trees — *a treebank* — count occurrences of each rule application and normalize wrt respective non-terminal
- Every PCFG built using a treebank is proper

# Computational Tasks

**Evaluation:** assessing the value of a given tree

⇒ multiply probabilities associated with each rule used in building the tree

(saw this already)

# Computational Tasks

**Generation:** producing sentences

⇒ monkeys again. Start with **S**, select derivation rule randomly according to the probability distribution, repeat for each resulting non-terminal until none left

If PCFG is proper, procedure will halt

# Computational Tasks

**Learning:** building a PCFG from a treebank

⇒ Count occurrences of each rule for  $X$ , divide result by number of all rules for  $X$

(saw this already)

# Probabilistic Inference

## Inference Tasks:

$$P(\textit{sentence})$$

(Marginalization)

$$P(\textit{tree}|\textit{sentence})$$

(Conditioning)

$$\textit{arg max}_{\textit{tree}} P(\textit{tree}|\textit{sentence})$$

(Completion)



# Probabilistic Inference

Consider marginalization

$$\begin{aligned} P(\textit{sentence}) &= P(w_1 w_2 \dots w_n \mid \mathbf{S}) \\ &= \sum_{\textit{tree} \in T} P(\textit{tree}) \end{aligned}$$

- Need to find the set  $T$  of all trees for *sentence*
- Need to compute each tree's probability
- Likely to lead to an exponential algorithm

For grammar  $\{S \rightarrow S S, S \rightarrow a\}$ , how many trees does  $a^n$  have?

# Efficient Probabilistic Inference

- Can use a version of CYK, an excellent CFG parser
- Only works on CFG in Chomsky Normal Form

# Efficient Probabilistic Inference

## Chomsky Normal Form

A CFG is in CNF if its every derivation rule has one of the two forms:

$$A \rightarrow BC$$

$$A \rightarrow w$$

where **A**, **B**, **C** are non-terminals and  $w$  is a terminal.

- Any CFG can be converted to CNF
- How to translate probabilities?  
estimate by sampling or calculate directly

# Efficient Probabilistic Inference

## Calculating CNF:

- Eliminate empty rules  $\mathbf{X} \rightarrow \epsilon$   
strike out RHS appearances of nullable terminals in all ways except one
- Eliminate unit rules  $\mathbf{X} \rightarrow \mathbf{Y}$   
if  $A$  derives  $B$  and  $B \rightarrow \phi$ , add  $A \rightarrow \phi$
- Eliminate terminals except singletons  
introduce new non-terminals as necessary
- Break rules  $\mathbf{X} \rightarrow \mathbf{Y}_1 \mathbf{Y}_2 \dots \mathbf{Y}_n, n > 2$   
introduce  $n - 2$  new non-terminals for each rule, replace rule with a set of new rules

# Efficient Probabilistic Inference

Our grammar in CNF:

S	→	NP VP	VP	→	V NP	N	→	time	V	→	like
NP	→	time	VP	→	V PP	N	→	arrow	V	→	flies
NP	→	N N	PP	→	P NP	N	→	flies	P	→	like

# Efficient Probabilistic Inference

## CYK — Cocke-Younger-Kasami Algorithm

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### Algorithm 1 CYK

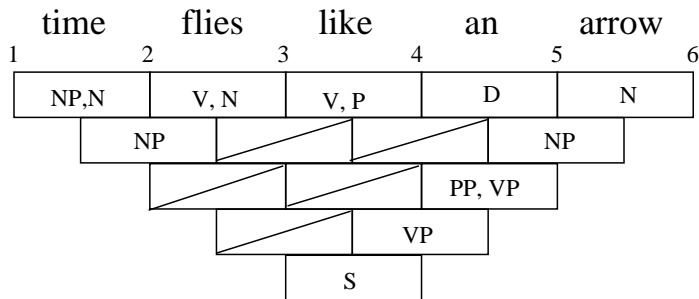
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**Require:** sentence =  $w_1 \dots w_n$ , and a CFG in CNF with nonterminals  $N^1 \dots N^m$ ,  
 $N^1$  is the start symbol

**Ensure:** parsed sentence

- 1: allocate matrix  $\beta \in \{0, 1\}^{n \times n \times m}$  and initialize all entries to 0
  - 2: **for**  $i \leftarrow 1$  to  $n$  **do**
  - 3:     **for all** rules  $N^k \rightarrow w_i$  **do**
  - 4:     |  $\beta[i, 1, k] \leftarrow 1$
  - 5: **for**  $j \leftarrow 2$  to  $n$  **do**
  - 6:     **for**  $i \leftarrow 1$  to  $n - j + 1$  **do**
  - 7:     | **for**  $l \leftarrow 1$  to  $j - 1$  **do**
  - 8:     | | **for all** rules  $N^k \rightarrow N^{k_1} N^{k_2}$  **do**
  - 9:     | | |  $\beta[i, j, k] \leftarrow \beta[i, j, k]$  OR  $(\beta[i, l, k_1]$  AND  $\beta[i + l, j - l, k_2])$
  - 10: **return**  $\beta[1, n, 1]$
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# Efficient Probabilistic Inference



Basic idea: exploiting CNF, build a hierarchical chart of non-terminals

# Efficient Probabilistic Inference

- Add probabilities to the mix  
NOT straight-forward, but has been done
- Obtain efficient marginalization for PCFGs



# Applications

- Monkeys
- Everything that needs to resolve ambiguity of NL parsing
- Detection of grammatical errors in text