Concurrent Object Oriented Languages Semaphores

wiki.eecs.yorku.ca/course/6490A

CSE 6490A

A semaphore is a datatype. Its values are nonnegative integers. A semaphore, say s, supports two atomic operations:

- V(s): increment s by 1.
- P(s): decrement s by 1 as soon as the result is nonnegative.

Using semaphores, we solve the following classical concurrency problems:

- The Producer-Consumer Problem (last lecture)
- The Readers-Writers Problem (already started last lecture)
- The Dining Philosophers Problem
- The Cigarette Smokers Problem

The readers and writers problem, due to Courtois, Heymans and Parnas, is another classical concurrency problem. It models access to a database. There are many competing threads wishing to read from and write to the database. It is acceptable to have multiple threads reading at the same time, but if one thread is writing then no other thread may either read or write. The problem is how do you program the reader and writer threads?

David Parnas

- Canadian early pioneer of software engineering
- Ph.D. from Carnegie Mellon University
- Taught at the University of North Carolina at Chapel Hill, the Technische Universität Darmstadt, the University of Victoria, Queen's University, McMaster University, and University of Limerick
- Won numerous awards including ACM SIGSOFT's "Outstanding Research" award



David Parnas

source: Hubert Baumeister

 Professor emeritus at the Catholic University of Leuven



Pierre-Jacques Courtois

source: www.info.ucl.ac.be/~courtois

CSE 6490A

The readers and writers share the following variable. semaphore mutex = 1; Reader: P(mutex); read; V(mutex); Writer: P(mutex); write; V(mutex);

Does it solve the readers-writers problem?



코 > - 코

Does it solve the readers-writers problem?

Answer

Yes!



ъ

<ロト < 同ト < 三ト

Does it solve the readers-writers problem?

Answer

Yes!

Question

Is it a satisfactory solution?



프 > 프

Does it solve the readers-writers problem?

Answer

Yes!

Question

Is it a satisfactory solution?

Answer

No!

CSE 6490A

< D > < P >

프 > 프

Why not?



æ –

イロト イロト イヨト イヨト

Why not?

Answer

It does not allow multiple readers to read at the same time.



<ロト < 同ト < 三ト

프 > 프

While a writer is writing, readers and writers arrive. Once the writer is done, can the readers start reading?

Options

While readers are reading, readers and writers arrive. Can the readers start reading?

While a writer is writing, readers and writers arrive. Once the writer is done, can the readers start reading?

Yes

Options

While readers are reading, readers and writers arrive. Can the readers start reading?

Yes

No reader is kept waiting unless a writer is writing.

While a writer is writing, readers and writers arrive. Once the writer is done, can the readers start reading?

Options

While readers are reading, readers and writers arrive. Can the readers start reading?

While a writer is writing, readers and writers arrive. Once the writer is done, can the readers start reading?

No

Options

While readers are reading, readers and writers arrive. Can the readers start reading?

No

If a writer wants to write, it writes as soon as possible.

In the dining philosophers problem, due to Dijkstra, five philosophers are seated around a round table. Each philosopher has a plate of spaghetti. The spaghetti is so slippery that a philosopher needs two forks to eat it. The layout of the table is as follows.



The life of a philosopher consists of alternative periods of eating and thinking. When philosophers get hungry, they try to pick up their left and right fork, one at a time, in either order. If successful in picking up both forks, the philosopher eats for a while, then puts down the forks and continues to think.

```
int N = 5;
philosopher(i):
while (true)
  think;
  takeForks(i);
  eat;
  putForks(i);
takeForks(i):
  . . .
putForks(i):
  . . .
```

프 > 프

```
int N = 5;
semaphore mutex = 1;
philosopher(i):
while (true)
  think;
  takeForks(i);
  eat;
  putForks(i);
takeForks(i):
P(mutex);
putForks(i):
V(mutex);
```

イロト イポト イヨト イヨトー

1

Question

Is this solution correct?



(日)

문▶ 문

Question

Is this solution correct?

Answer

Yes!



æ

Question

Is this solution correct?

Answer

Yes!

Question

Does this solution allow two philosophers to eat at the same time?

<ロト < 同ト < 三ト

프 > 프

Question

Is this solution correct?

Answer

Yes!

Question

Does this solution allow two philosophers to eat at the same time?

Answer

No!



 $\langle \Box \rangle \langle \Box \rangle$

프 > 프

In the Cigarette Smokers Problem, due to Patil, there are an agent and three smokers. The smokers loop forever, first waiting for ingredients, then making and smoking cigarettes. The ingredients are tobacco, paper, and matches. We assume that the agent has an infinite supply of all three ingredients, and each smoker has an infinite supply of one of the ingredients; that is, one smoker has matches, another has paper, and the third has tobacco.

The agent repeatedly chooses two different ingredients at random and makes them available to the smokers. Depending on which ingredients are chosen, the smoker with the complementary ingredient should pick up both resources and proceed. For example, if the agent puts out tobacco and paper, the smoker with the matches should pick up both ingredients, make a cigarette, and then signal the agent.

・ロト ・ 雪 ト ・ ヨ ト ・

Question

Can we implement the random choices made by the agent using concurrency?



코 > - 코

Question

Can we implement the random choices made by the agent using concurrency?

Answer

Yes, we can introduce a thread for each choice and run them concurrently.



Question

Can we implement the random choices made by the agent using concurrency?

Answer

Yes, we can introduce a thread for each choice and run them concurrently.

Question

How can we ensure that at most one of those threads executes at any time?

Can we implement the random choices made by the agent using concurrency?

Answer

Yes, we can introduce a thread for each choice and run them concurrently.

Question

How can we ensure that at most one of those threads executes at any time?

Answer

Introduce a binary semaphore.

イロト イポト イヨト イヨト

Problem

Use semaphores paper and tabacco for the agent who puts out tobacco and paper, signal the smoker with the matches.



Shared variables

semaphore agent = 1 semaphore match = 0 semaphore paper = 0 semaphore tabacco = 0

Agent (paper and tabacco):

```
P(agent); V(tabacco); V(paper)
```

Agent (paper and matches):

```
P(agent); V(paper); V(match)
```

Agent (tabacco and matches):

```
P(agent); V(tabacco); V(match)
```

▲□ ▶ ▲ ■ ▶ ▲ ■ ▶ ● ● ● ● ● ●

Smoker (paper):

P(tabacco); P(match); V(agent)

Smoker (matches):

```
P(tabacco); P(paper); V(agent)
```

Smoker (tabacco):

P(paper); P(match); V(agent)

э.

Is this solution correct?



æ

<ロト < 同ト < 三ト

Is this solution correct?

Answer

No, because it can deadlock.



<ロト < 同ト < 三ト

포 > 프