Smaller Models EECS 4315

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```
public class Counter extends Thread
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  private int value;
  public Counter()
    this.value = 0;
  }
  . . .
 }
```

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Counter Class

```
public void run()
{
   this.value++;
}
```

- 0: aload_0
- 1: dup
- 2: getfield
- 5: iconst_1
- 6: iadd
- 7: putfield
- 10: return

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```
public class Main
ł
  public static void main(String[] args)
  ł
    Counter one = new Counter();
    Counter two = new Counter();
    one.start();
    two.start();
ł
```

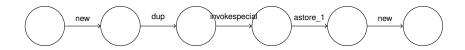
0:	new	11:	dup	20:	aload_2
3:	dup	12:	invokespecial	21:	invokevirt
4:	invokespecial	15:	astore_2	24:	return
7:	astore_1	16:	aload_1		
8:	new	17:	invokevirtual		
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Draw the corresponding state-transition diagram.

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State-Transition Diagram



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State-Transition Diagram



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Combine the first ten transitions into one.

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Combine the first ten transitions into one.



The actions of the labelled transition system are sequences of bytecode instructions.

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State-Transition Diagram



Next instructions for the main thread:

- 20: aload_2
- 21: invokevirtual
- 24: return

Next instructions for the thread one:

- 0: aload_0
- 1: dup
- 2: getfield
- 5: iconst_1
- 6: iadd
- 7: putfield

EECS 4315

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Can the bytecode instructions corresponding to the **run** invocation be modelled as a single transition?

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Answer

Yes.

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Answer

Yes.

Question

Why?

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Can the bytecode instructions corresponding to the **run** invocation be modelled as a single transition?

Answer

Yes.

Question

Why?

Answer

Because the execution of this method does not impact the other threads.

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Combining Bytecode Instructions

- We combine the first ten bytecode instructions since there is only one thread.
- We combine the bytecode instructions corresponding to the **run** invocation because those do not impact the other threads.

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Combining Bytecode Instructions

- We combine the first ten bytecode instructions since there is only one thread.
- We combine the bytecode instructions corresponding to the **run** invocation because those do not impact the other threads.

General idea

Combine those bytecode instructions that do not impact other threads.

Given all the bytecode instructions, determine for a specific instruction whether it impacts other threads.

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Given all the bytecode instructions, determine for a specific instruction whether it impacts other threads.

Question

Give an algorithm that solves the problem.

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Question

Give an algorithm that solves the problem.

Question

Impossible!

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Which other problems cannot be solved?

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Which other problems cannot be solved?

Answer

The halting problem: given code and input for that code, determine whether the code terminates.

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Given all the bytecode instructions, determine for a specific instruction whether it impacts other threads.

Question

Prove that the problem cannot be solved.

General idea

Combine those bytecode instructions for which we can prove that they do not impact other threads.

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General idea

Combine those bytecode instructions for which we can prove that they do not impact other threads.

The idea of combining consecutive transitions labelled with invisible (outside the current thread) actions into a single transition is due to Patrice Godefroid.

General idea

Combine those bytecode instructions for which we can prove that they do not impact other threads.

The idea of combining consecutive transitions labelled with invisible (outside the current thread) actions into a single transition is due to Patrice Godefroid.

Examples of invisible actions

- Reading or writing an attribute that can be proved to be not shared.
- Reading or writing a local variable.
- . . .

Patrice Godefroid

- Ph.D. degree in Computer Science from the University of Liege, Belgium
- Worked at Bell Laboratories.
- Currently at Microsoft Research.



Source: Patrice Godefroid

Java PathFinder can generate state space diagrams in the form of dot files.

Target=Main
classpath=.
listener=gov.nasa.jpf.listener.SimpleDot

Java PathFinder can generate state space diagrams in the form of dot files.

Target=Main
classpath=.
listener=gov.nasa.jpf.listener.StateSpaceDot

A semaphore is a nonnegative integer, say s, with two atomic operations:

- V(s): increment s by 1.
- P(s): decrement s by 1 as soon as the result is nonnegative.

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- P(s): decrement s by 1 as soon as the result is nonnegative.

The Java package java.util.concurrent contains the class Semaphore. The operations V and P are represented by the methods release and acquire.