

Linear Temporal Logic

EECS 4315

www.eecs.yorku.ca/course/4315/

Linear temporal logic (LTL) is a logic to reason about systems with nondeterminism.

The logic was introduced by Amir Pnueli.

A. Pnueli. The temporal logic of programs. In *Proceedings of the 18th IEEE Symposium on Foundations of Computer Science*, pages 46–67. Providence, RI, USA, October/November 1977. IEEE.

Amir Pnueli (1941–2009)

- Recipient of the Turing Award (1996)
- Recipient of the Israel prize (2000)
- Foreign Associate of the U.S. National Academy of Engineering (1999)
- Fellow of the Association for Computing Machinery (2007)



Source: David Monniaux

Definition

LTL is defined by the following grammar.

$$f ::= a \mid f \wedge f \mid \neg f \mid \bigcirc f \mid f \mathbf{U} f$$

where a is an atomic proposition.

An atomic proposition represents a basic property (such as the value of a particular variable being even).

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Does the execution path



satisfy the atomic proposition **green**?

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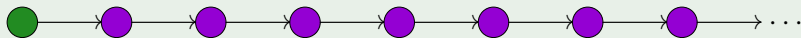
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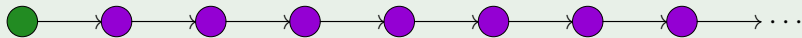
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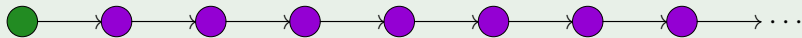
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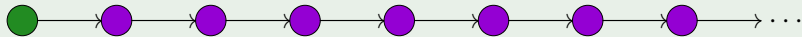
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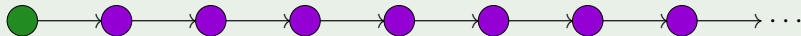
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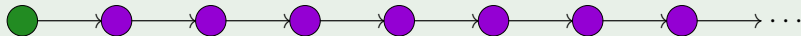
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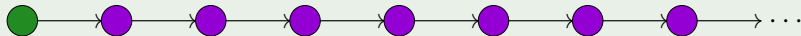
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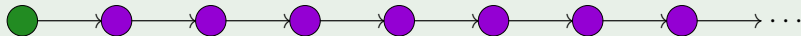
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Answer

Yes.

As usual

$$\begin{aligned}\text{true} &= a \vee \neg a \\ f_1 \vee f_2 &= \neg(\neg f_1 \wedge \neg f_2) \\ f_1 \Rightarrow f_2 &= \neg f_1 \vee f_2\end{aligned}$$

Also

$$\begin{aligned}\diamond f &= \text{true} \text{ U } f \quad (\text{eventually } f) \\ \square f &= \neg \diamond \neg f \quad (\text{always } f)\end{aligned}$$

Xf : $\bigcirc f$
 Ff : $\diamond f$
 Gf : $\square f$

State Space Diagram

Question

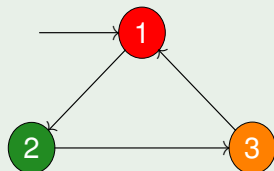
Draw the state space diagram of a model of a traffic light.

State Space Diagram

Question

Draw the state space diagram of a model of a traffic light.

Answer



Note: the transitions are not labelled, but the states are labelled.

Definition

A transition system is a tuple $\langle S, L, I, \rightarrow, \ell \rangle$ consisting of

- a set S of states,
- a set L of labels,
- A set $I \subseteq S$ of initial states,
- a transition relation $\rightarrow \subseteq S \times S$ such that for all $s \in S$ there exists a $s' \in S$ such that $s \rightarrow s'$, and
- a labelling function $\ell : S \rightarrow 2^L$.

2^L denotes the set of subsets of L .

Question

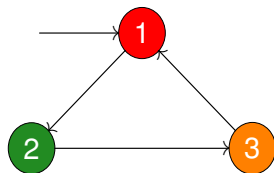
What is $2^{\{1,2,3\}}$?

Question

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Answer

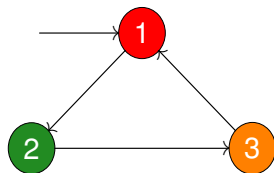
$\{\emptyset, \{1\}, \{2\}, \{3\}, \{1, 2\}, \{1, 3\}, \{2, 3\}, \{1, 2, 3\}\}$.



Question

Give the transition system modelling a traffic light.

Transition System

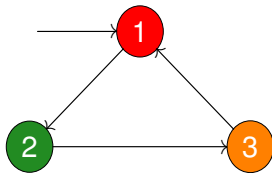


Question

Give the transition system modelling a traffic light.

Answer

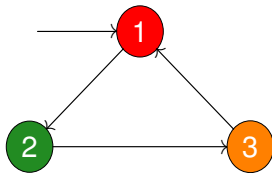
$$\langle \{1, 2, 3\}, \{\text{red}, \text{green}, \text{orange}\}, \\ \{1\}, \{1 \rightarrow 2, 2 \rightarrow 3, 3 \rightarrow 1\} \\ \{1 \mapsto \{\text{red}\}, 2 \mapsto \{\text{green}\}, 3 \mapsto \{\text{orange}\}\} \rangle$$



$$f ::= a \mid f \wedge f \mid \neg f \mid \bigcirc f \mid f \mathbf{U} f$$

Question

Which LTL formula expresses
“initially the light is red.”



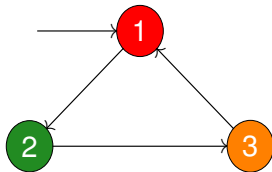
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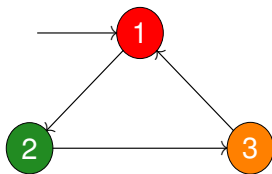
The LTL formula **red**.



$$f ::= a \mid f \wedge f \mid \neg f \mid \bigcirc f \mid f U f$$

Question

Which LTL formula expresses
“initially the light is red and next it becomes green.”



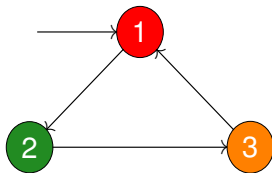
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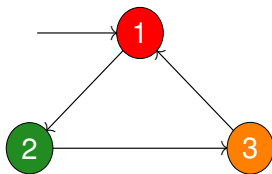
The LTL formula **red** \wedge \bigcirc **green**



$$f ::= a \mid f \wedge f \mid \neg f \mid \bigcirc f \mid f \mathbf{U} f$$

Question

Which LTL formula expresses
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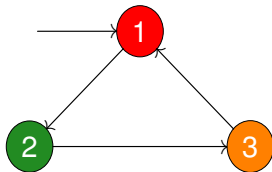
$$f ::= a \mid f \wedge f \mid \neg f \mid \bigcirc f \mid f \text{ U } f$$

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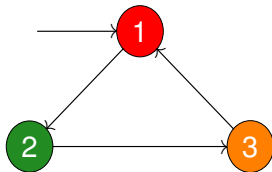
The LTL formula true U orange = \diamond orange



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Which LTL formula expresses
“the light is infinitely often red.”



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“the light is infinitely often red.”

Answer

The LTL formula $\square \diamond \text{red}$

Question

What does the LTL formula $\Box(\text{green} \Rightarrow \neg \bigcirc \text{red})$ express?

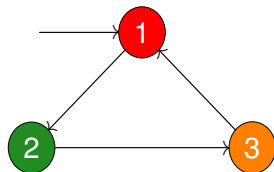
Question

What does the LTL formula $\Box(\text{green} \Rightarrow \neg\bigcirc\text{red})$ express?

Answer

“Once green, the light cannot become red immediately”

Execution Paths



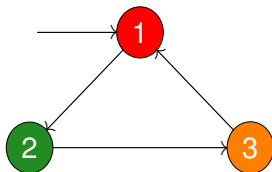
Definition

$Paths(s)$ is the set of path starting in state s .

Question

What is $Paths(2)$?

Execution Paths



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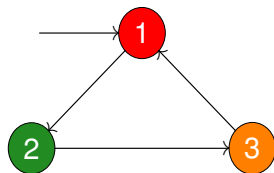
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Answer

$\{231231231231 \dots\}$

Execution Paths



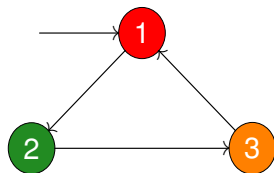
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Let $p \in Paths(s)$ and $n \geq 0$. Then $p[n]$ is the $(n + 1)^{th}$ state of the path p .

Question

Let $p = 123123 \dots$. What is $p[3]$?

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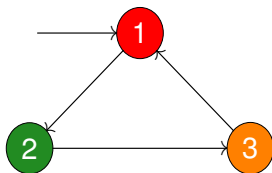
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Execution Paths



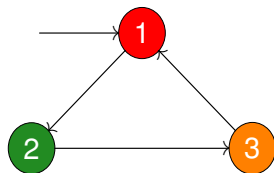
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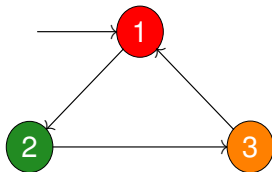
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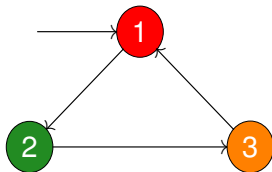
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312312...



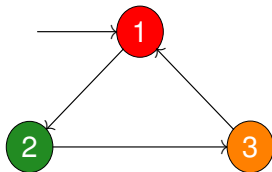
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$123123 \dots \models \text{green?}$



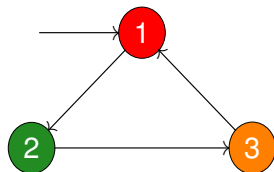
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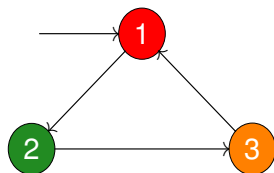
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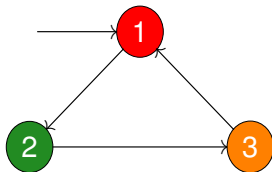
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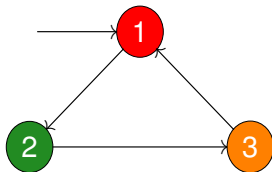
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$123123 \dots \models \text{red} \wedge \bigcirc \text{green}?$



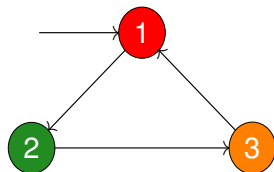
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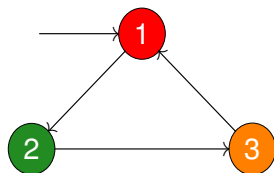
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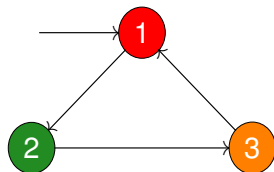
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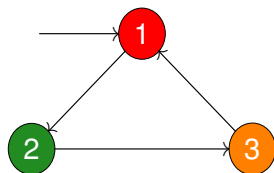
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Answer

Yes.

Definition

$$\begin{aligned} p \models a & \text{ iff } a \in \ell(p[0]) \\ p \models f \wedge \psi & \text{ iff } p \models f \wedge p \models \psi \\ p \models \neg f & \text{ iff } p \not\models f \\ p \models \bigcirc f & \text{ iff } p[1..] \models f \\ p \models f \cup \psi & \text{ iff } \exists i \geq 0 : p[i..] \models \psi \wedge \forall 0 \leq j < i : p[j..] \models f \end{aligned}$$

$$TS \models f \text{ iff } \forall s \in I : s \models f$$

where

$$s \models f \text{ iff } \forall p \in \text{Paths}(s) : p \models f$$

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